

Lazy interactions – back to the future

Simon Thompson, University of Kent

`System.IO.interact :: (String -> String) -> IO ()`

```
System.IO.interact :: (String -> String) -> IO ()
```

```
interact f = do s <- getContents  
                putStr (f s)
```

$f :: (\text{Input} \rightarrow \text{Output})$

The output of the program is a function of its input.

$f :: (\text{Input} \rightarrow \text{Output})$

3

23

45

67



(2, 23)

(1, 68)

(0, 135)

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Interaction

=

input / output
interleaving

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(0, 135)

Interaction

=

input / output
interleaving

Interleaving
determined
by lazy evaluation

The essence of laziness

f ⊥
= "type now" ++ ⊥

f ("echo" ++ ⊥)
= ... ++ "ohce" ++ ⊥

The essence of laziness

```
f ⊥  
= "type now" ++ ⊥  
  
f ("echo" ++ ⊥)  
= ... ++ "ohce" ++ ⊥
```

Lazy interactions are determined by the behaviour of the function on **partial data**.

Demo

“Seat of the pants?”

```
necho ys  
  = "Prompt: " ++ [head ys] ++ "\n" ++ necho (tail ys)
```

vs

```
necho (x:xs)  
  = "Prompt: " ++ [x] ++ "\n" ++ necho xs
```

“Seat of the pants?”

```
necho ~(x:xs)  
  = "Prompt: " ++ [x] ++ "\n" ++ necho xs
```

vs

```
necho (x:xs)  
  = "Prompt: " ++ [x] ++ "\n" ++ necho xs
```

“Seat of the pants?”

nec

=

vs

nec

=

Let's build a model
of interactions and
how to combine
them together ...

Back to the future?

YEAR OF PROGRAMMING

The 1987 University of Texas Year of Programming was established early in 1986, in response to a proposal by Profs. J. C. Browne and J. Misra, with the following goals:

- 1) to advance the art and science of programming by bringing leading scientists together for discussions and collaboration;
- 2) to disseminate among leading practitioners the best of what has been learned about the theory and practice of programming; and

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The tutorial, which provided an introduction to lazy functional programming, consisted of lectures interspersed with programming sessions (conducted with pencil and paper) attended by the lecturers and several teaching assistants. Major topics included data types, polymorphism, recursion and induction, lists, domain theory, program synthesis, and several case studies.

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This institute elicited particular enthusiasm among a group of UT graduate students, who circulated among themselves, and subsequently presented to the UT Department of Computer Sciences, a petition calling on the department "to make Functional Programming a more visible priority in the department... [through] recruitment of faculty engaged in research in the field [and] more formal contacts with private research and other departments...".

Research Topics in Functional Programming

EDITED BY DAVID A. TURNER

UNIVERSITY OF TEXAS AT AUSTIN YEAR OF PROGRAMMING SERIES

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response to a propos

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Back to the future?

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The power of retrospection ...

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... how we can bring it up to date?

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Transliterating
from Miranda
to Haskell

Back to the future?

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The power of retrospection ...

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Transliterating
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Building a
formal model of
interactions, with
some proofs ...

$(\text{Input}, a) \rightarrow (\text{Input}, b, \text{Output})$

Functions with IO side effects

Build by composition

(Input, a) -> (Input, b, Output)

Functions with IO side effects

Build by composition

Interactions with states

State changes type between steps ...

... can add, remove, and modify what's there.

(Input, a) -> (Input, b, Output)

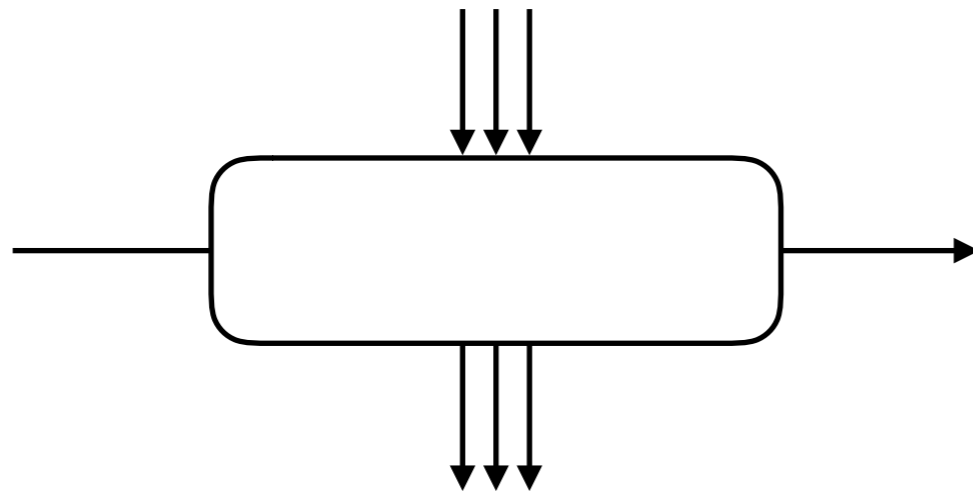
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Basic types ...

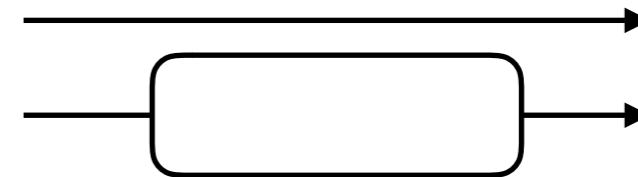
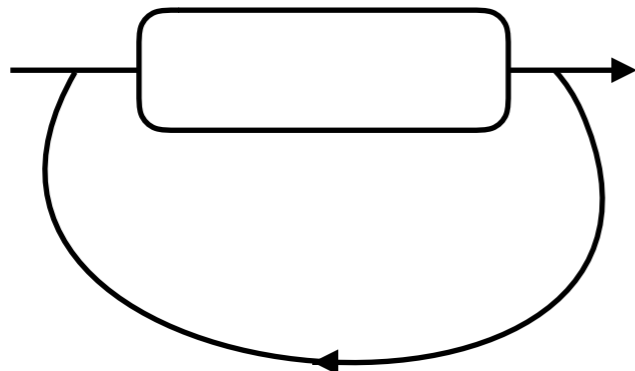
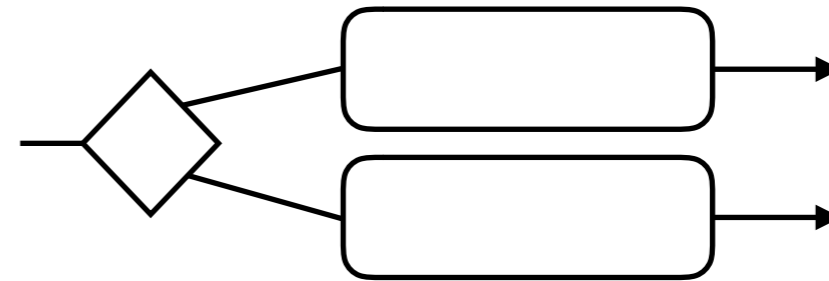
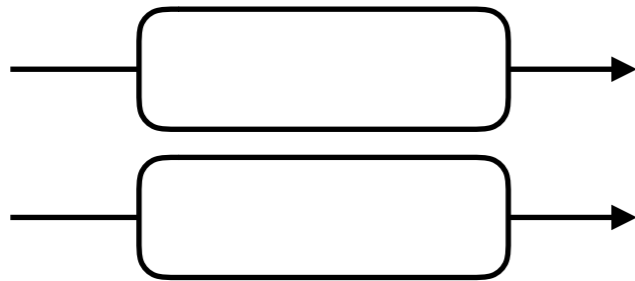
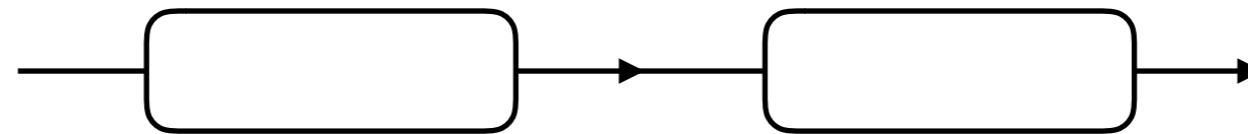
```
type Interact a b
  = (Input, a) -> (Input, b, Output)
```

```
type Condition a
  = (Input, a) -> Bool
```

```
type Input  = [String]
type Output = [String]
```

How do we put these together?

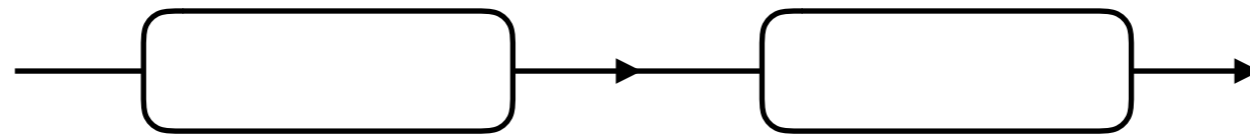
We assume that all diagrams are well-typed



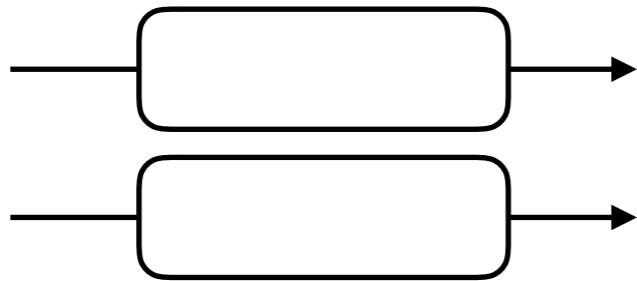
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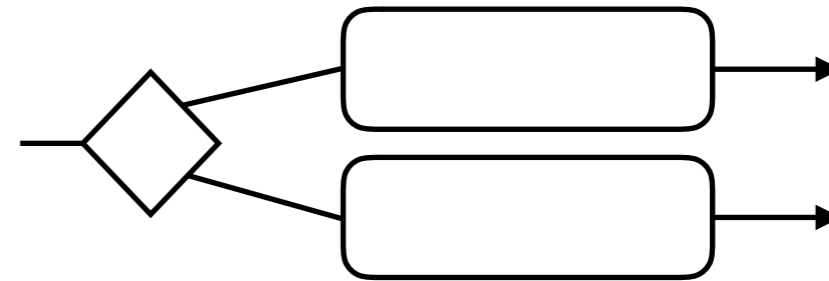
sq



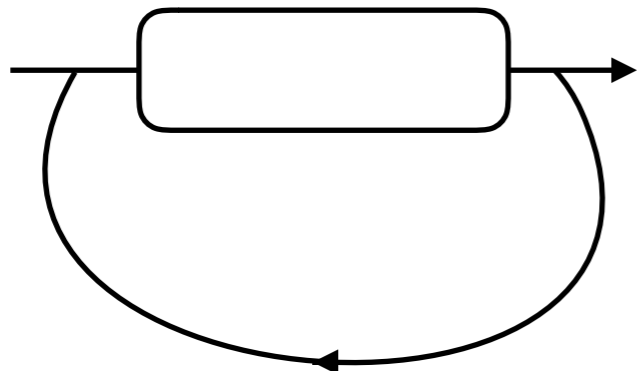
par



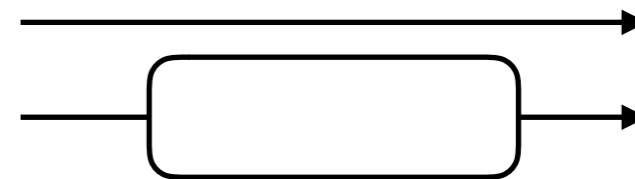
alt



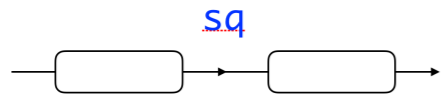
while



pass_on_l



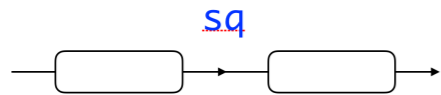
Sequencing ... key combinator



`sq :: Interact a b -> Interact b c -> Interact a c`

```
sq inter1 inter2 x
  = make_Output out1 (inter2 (rest,st))
  where (rest,st,out1) = inter1 x
```


Sequencing ... key combinator



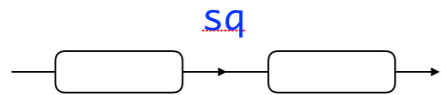
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`make_Output :: Output -> (Input,a,Output) -> (Input,a,Output)`

```
make_Output piece (input,st,out) = (input,st,piece++out)
```

Sequencing ... key combinator



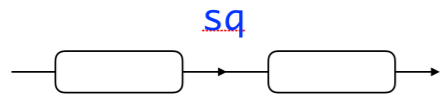
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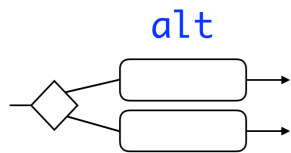
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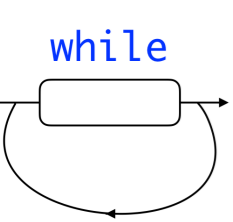
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Alternation and repetition



$\text{alt} :: \text{Cond } a \rightarrow \text{Interact } a \ b \rightarrow \text{Interact } a \ b \rightarrow \text{Interact } a \ b$

```
alt cond inter1 inter2 x
  | cond x      = inter1 x
  | otherwise   = inter2 x
```



$\text{while} :: \text{Cond } a \rightarrow \text{Interact } a \ a \rightarrow \text{Interact } a \ a$

```
while cond inter
  = whi
  where
    whi = alt cond (inter `sq` whi) null
```

“Passing parameters”

```
pass_param :: Interact a b ->  
            (b -> Interact () d) ->  
            Interact a d
```

```
pass_param int f (input,st)  
  = (rest,final,out1++out)  
  where  
    ~(inter1,st1,out1) = int (input,st)  
    ~(rest,final,out)  = (f st1) (inter1,())
```

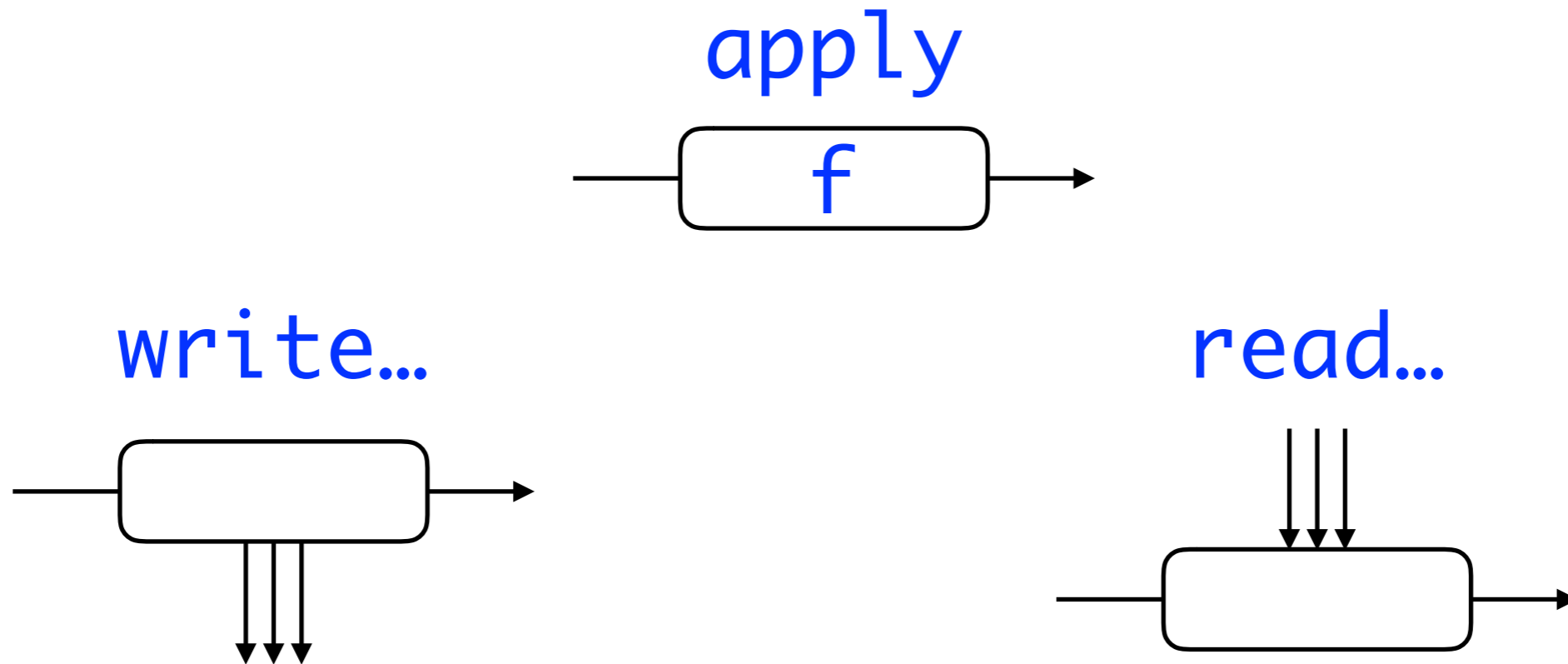
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```

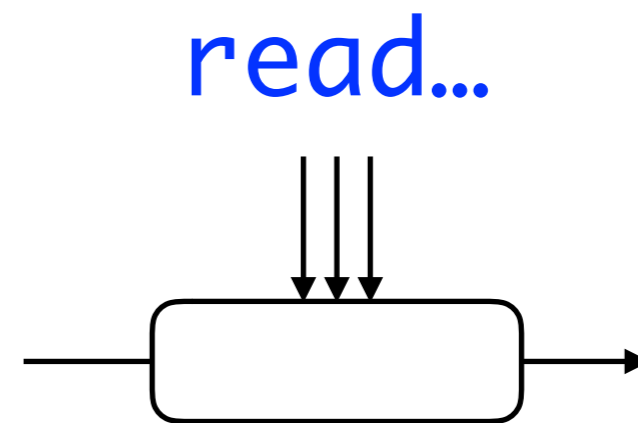
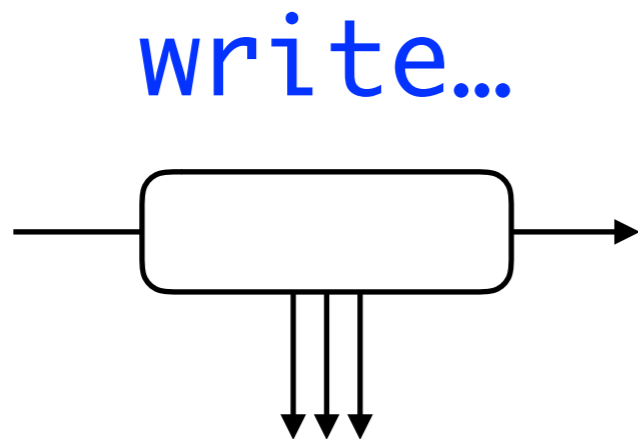
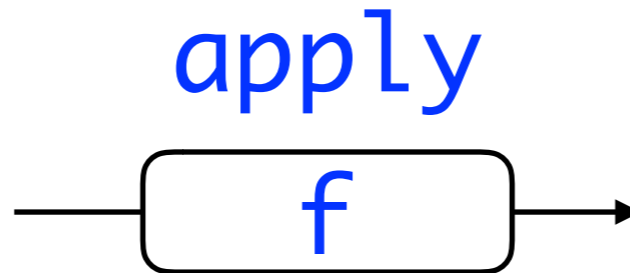
And some primitives ...

We assume that
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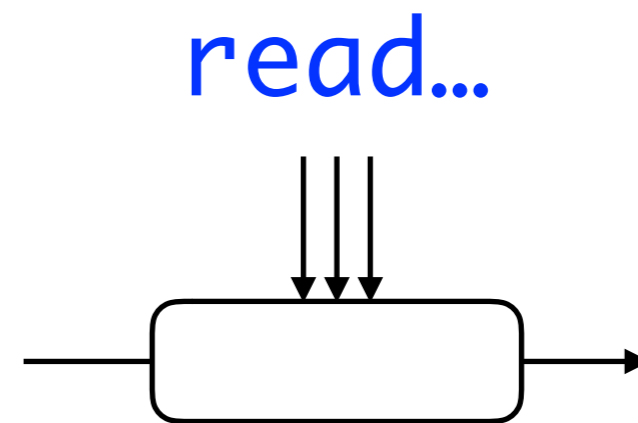
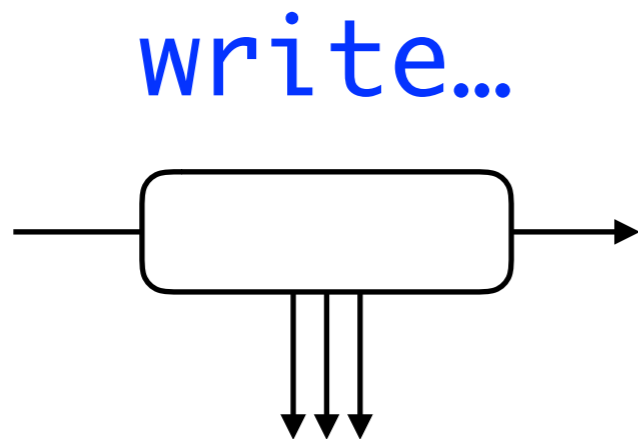
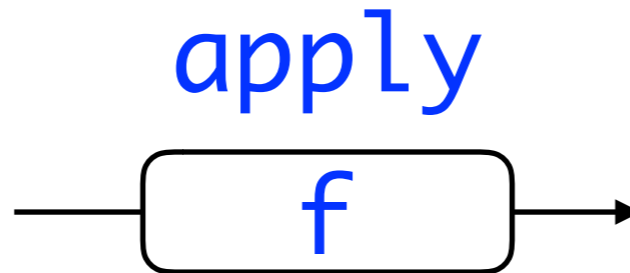
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forget start change wait

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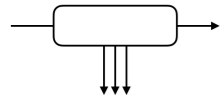


forget start change wait

run :: Interact a b -> a -> IO ()

And some primitives ...

write...



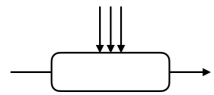
`write :: String -> Interact a a`

`write outstring (input,st)
= (input,st,[outstring])`

`run :: Interact a b -> a -> IO ()`

`run inter st
= interact (\chs ->
case inter (split chs,st) of
(_,_,out) -> join out ++ "\n")`

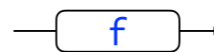
read...



`readin :: Interact () String`

`readin (input,())
= (tail input, head input,[])`

apply



`apply :: (a -> b) -> Interact a b`

`apply f (input,st)
= (input, f st, [])`

Demo

Copy input

```
copy :: Interact () ()
```

```
copy = while (\_ -> True) (readin `sq` writeout id)
```

Copy input

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copy :: Interact () ()
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copy = readin `sq` writeout id `sq` copy
```

Copy input

A little
meta-
circularity

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copy :: Interact () ()
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copy = while (\_ -> True) (readin `sq` writeout id)
```

```
copy :: Interact () ()
```

```
copy = readin `sq` writeout id `sq` copy
```

Input N then sum N numbers

```
collector :: Interact () (Int,Int)
```

```
collector
```

```
=  getInt `sq`  
   add_val_right 0 `sq`  
   while ((>(0::Int)).fst.snd)  
     (add_val_left () `sq`  
      pass_on getInt `sq`  
      apply (\(p,(m,s))->(m-1,s+p)) `sq`  
      wait `sq`  
      showkeep)
```

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```

counter

Input N then sum N numbers

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      wait `sq`  
      showkeep)
```

```
counter  
(counter,sum)
```

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```

```
counter  
(counter,sum)  
(counter,sum)
```

Input N then sum N numbers

```
collector :: Interact () (Int,Int)
```

```
collector
```

```
  = getInt `sq`                                counter  
    add_val_right 0 `sq`                       (counter,sum)  
    while ((>(0::Int)).fst.snd)                (counter,sum)  
      (add_val_left () `sq`                    ((), (counter,sum))  
      pass_on getInt `sq`  
      apply (\(p,(m,s))->(m-1,s+p)) `sq`  
      wait `sq`  
      showkeep)
```

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   while ((>(0::Int)).fst.snd)                (counter,sum)  
     (add_val_left () `sq`                    ((), (counter,sum))  
     pass_on getInt `sq`                     (Int,(counter,sum))  
     apply (\(p,(m,s))->(m-1,s+p)) `sq`  
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```

Input N then sum N numbers

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collector :: Interact () (Int,Int)
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```
=  getInt `sq`                                counter  
   add_val_right 0 `sq`                        (counter,sum)  
   while ((>(0::Int)).fst.snd)                 (counter,sum)  
     (add_val_left () `sq`                       ((), (counter,sum))  
     pass_on getInt `sq`                         (Int,(counter,sum))  
     apply (\(p,(m,s))->(m-1,s+p)) `sq`         (counter,sum)  
     wait `sq`  
     showkeep)
```


Input N then sum N numbers

Make the state abstract, with accessors, mutators etc.

```
collector :: Interact () (Int,Int)
```

```
collector
```

```
= getInt `sq`
  add_val_right 0 `sq`
  while ((>(0::Int)).fst.snd)
    (add_val_left () `sq`
     pass_on getInt `sq`
     apply (\(p,(m,s))->(m-1,s+p)) `sq`
     wait `sq`
     showkeep)
      counter
      (counter,sum)
      (counter,sum)
      ((), (counter,sum))
      (Int,(counter,sum))
      (counter,sum)
      :-)
```

Input N then sum N numbers

```
collectNums :: Interact Int Int
```

```
collectNums  
  = addNum  
    `pass_param`  
      (\n -> start 0 `sq`  
              seqList (replicate n addNum) `sq`  
              write "finished")
```


Input N then sum N numbers

Leave the
internal state
and synthesise
a program.

```
collectNums :: Interact Int Int
```

```
collectNums
```

```
  = addNum
```

```
    `pass_param`
```

```
      (\n -> start 0 `sq`
```

```
        seqlist (replicate n addNum) `sq`
```

```
        write "finished")
```

Looking back

All the ingredients were there ...

Higher-order functions

Lazy evaluation

Pattern matching

Algebraic data types

... well, almost all

Miranda had no `lambda`, or `let`.

- A variant of “point-free” style:
the need to name abstractions.

Equality overloaded ... but not `show`, ...

Few established “design patterns”

The model mixes aspects of

- Monad
- Arrow
- Applicative

The linguistic turn ...

Can see this as a *shallow embedding* of an interaction language.

What would happen if we made that deep?

The linguistic turn ...

Can see this as a *shallow embedding* of an interaction language.

What would happen if we made that deep?

```
data Inter =  
  While Cond Inter I  
  Alt Cond Inter Inter I  
  Seq Inter Inter I  
  ...  
  
interpret ::  
  Inter -> Interact Int Int
```

The linguistic turn ...

Can see this as a *shallow embedding* of an interaction language.

What would happen if we made that deep?

Questions of reflection, dependent types etc.

```
data Inter =  
  While Cond Inter I  
  Alt Cond Inter Inter I  
  Seq Inter Inter I  
  ...  
  
interpret ::  
  Inter -> Interact Int Int
```


Types

The fundamental scope of values hasn't changed ...

... but their classifications have.

Roles for e.g. GADTs, dependency here, especially with DSLs?

Time to look at Fudgets again?

α F, Fudget, et al



The Fudget type

Types

```
data F a b = F (FSP a b)
  instance FudgetIO F
  instance StreamProcIO F
type Fudget a b = F a b
type FSP a b = SP (FEvent a) (FCommand b)
type TEvent = (Path, FResponse)
type TCommand = (Path, FRequest)
type FEvent a = Message TEvent a
type FCommand a = Message TCommand a
```

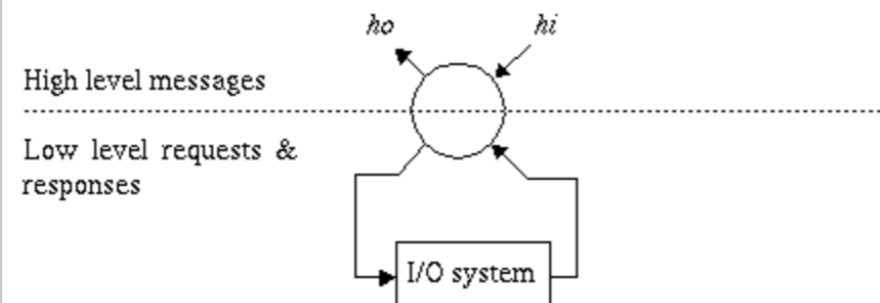
```
data SP a b
```

```
data Message a b = Low a | High b
```

Description

A *fudget* is a stream processor with high level streams and low level streams. The high level streams are used for communication between fudgets within a program. The low level streams are for communication with the I/O system.

$F\ hi\ ho$ is the Fudget type. hi is the type of high level input messages and ho is the type of high level output messages.



<http://www.altocumulus.org/Fudgets/>

Compilation

Libraries

Interop e.g FFI

APIs

???

Tools

Concurrency

Community

And what hasn't happened?

Routine verification ... semantics.

Compilers derived from semantics.

The end of the program as text.

Special purpose parallel hardware.

<https://github.com/simonjohnthompson/Interaction>