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3 Authorisation in Grid computing

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7
8 **Abstract** This paper briefly surveys how authorisation in Grid computing has
9 evolved during the last few years, and presents the latest developments in which
10 Grid applications can utilise a policy controlled authorisation infrastructure to
11 make decisions about which users are allowed to perform which actions on which
12 Grid resources. The paper describes the Global Grid Forum SAML interface for
13 connecting policy based authorisation infrastructures to Grid applications, and then
14 describes the PERMIS authorisation infrastructure which has implemented this
15 interface. The paper concludes with suggestions about how this work will evolve in
the future.

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17 Introduction

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21 Grid computing allows resources, including large
22 scale expensive ones such as genetic databases, to
23 be shared between members of virtual organisa-
24 tions (VOs). However, if an organisation is to allow
25 its resources to be shared amongst its VO partners,
26 it needs to be able to determine who is authorised
27 to access these resources in which ways, and who
28 is not. Ideally each resource owner should be able
29 to set the policy determining the rules for who is
30 authorised to do what, and then leave the author-
31 isation infrastructure to enforce this policy. The
32 resource owner should not be expected to

individually identify and name each VO user who 36
is to access his/her resource, as this soon becomes 37
unwieldy and costly to manage. The resource 38
owner should be able to delegate to other partners 39
in the VO the ability to identify and nominate the 40
users from their respective domain who are to be 41
allowed to use his/her resource, leaving him/her 42
to simply determine what type of access to grant 43
to the different categories of user. It is only very 44
recently that we have been able to achieve this, as 45
will be described here. 46
47
48

The rest of this paper is structured as follows. 49
Next section provides a brief history of Grid 50
authorisation. Then the current Global Grid Forum 51
draft authorisation interface is described. Further 52
the PERMIS policy based authorisation infrastruc- 53
ture that is compatible with the GGF interface is 54
described. Finally last section concludes and looks 55
at possible future work in this area. 56
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2 A brief history of Grid authorisation

3 One of the earliest attempts at providing authorisation in VOs was in the form of the Globus Toolkit
4 Gridmap files (Sotomayor). This file simply holds
5 a list of the authenticated distinguished names of
6 the Grid users and the equivalent local user
7 account names that they are to be mapped into.
8 Access control to a resource is then left up to the
9 local operating system and application access
10 control mechanisms. As can be seen, this neither
11 allows the local resource administrator to set
12 a policy for who is allowed to do what, nor does
13 it minimise his/her workload. On the contrary it
14 maximises the work of the resource administrator
15 since (s)he must first pre-configure the Grid appli-
16 cation with the names of every VO user who is to
17 be allowed to access the Grid resource, and then
18 (s)he must set the access controls on the local
19 operating system and/or application to ensure that
20 the local user names are restricted in what they
21 are allowed to do with the resource. The system is
22 neither scalable nor flexible, and there is no way
23 to distribute the administrative task throughout
24 the VO. Consequently several ways were devel-
25 oped to improve upon this.

27 The Community Authorisation Service (CAS)
28 (Pearlman et al., 2002) was the next attempt by
29 the Globus team to improve upon the manage-
30 ability of user authorisation. CAS allows a resource
31 owner to grant access to a portion of his/her
32 resource to a VO (or community – hence the name
33 CAS), and then let the community determine who
34 can use this allocation. The resource owner thus
35 partially delegates the allocation of authorisation
36 rights to the community. This is achieved by having
37 a CAS server, which acts as a trusted intermediary
38 between VO users and resources. Users first con-
39 tact the CAS asking for permission to use a Grid
40 resource. The CAS consults its policy (which
41 specifies who has permission to do what on which
42 resources) and if granted, returns a digitally self-
43 signed capability to the user optionally containing
44 policy details about what the user is allowed to do
45 (as an opaque string). The user then contacts the
46 resource and presents this capability. The resource
47 checks that the capability is signed by a known and
48 trusted CAS and if so maps the CAS's distinguished
49 name into a local user account name via the
50 Gridmap file. Consequently the Gridmap file now
51 only needs to contain the name of the trusted CAS
52 servers and not all the VO users. This substantially
53 reduces the work of the resource administrator.
54 Further, determining who should be granted capa-
55 bilities by the CAS server is the task of other

managers in the VO community, so this again
relieves the burden of resource managers. For
finer grained access control, the resource can
additionally call a further routine, passing to it
the opaque policy string from the capability, and
using the returned value to refine the access rights
of the user. Unfortunately this part of the CAS
implementation (policy definition and evaluation
routine) was never fully explored and developed
by the Globus team. Research by other groups into
policy controlled authorisation infrastructures
overtook the CAS work.

European researchers were never content with
the capabilities of either the manually generated
Gridmap file or the CAS. Consequently the EU
DataGrid and DataTAG projects developed the
Virtual Organisation Membership Service (VOMS)
(Alfieri et al., 2003) as a way of delegating the
authorisation of users to managers in the VO. VOMS
has gone through a number of iterations in its
development. Initially it was a system for dynam-
ically creating Gridmap files from LDAP directories
containing details about VO users. Resources could
pull a Gridmap file from this periodically. Thus the
resource owner never had to actually create or
manage the Gridmap file. This system, however,
was not scalable. The EU work then evolved into
a push system in which the VOMS server digitally
signed a “pseudo-certificate” for the VO user to
present to the resource. This pseudo-certificate
could contain a local user account name, in which
case no Gridmap file would be needed, or it could
contain other privileges or group membership de-
tails, in which case software would be needed by
the resource to interpret this information and
grant appropriate rights. The software they devel-
oped for this is called the Local Centre Author-
isation Service (LCAS) (Steenbakkers, 2003). LCAS
makes its authorisation decision based upon the
user's certificate and the job specification, which
is written in job description language (JDL) format.

In its current re-incarnation, the VOMS server
now produces short-lived X.509 attribute certifi-
cates (ACs) (ISO 9594-8/ITU Rec. X.509, 2001) for
the user to push to the resource. This design is
similar in concept to the CAS, but differs in
message format and syntax. In VOMS the name of
the user is presented to the resource instead of the
name of the CAS server, and user attributes are
presented instead of opaque policy statements.
The message construct is signed by the VOMS server
instead of the CAS server, and is a standard X.509
AC instead of a proprietary capability. However,
what neither VOMS nor CAS nor LCAS provides is the
ability for the resource administrator to set the

111 policy for access to his/her resource and then let
 112 the authorisation infrastructure enforce this policy
 113 on his/her behalf. This is what systems like Akenti
 114 (Johnston et al., 1998), PERMIS (Chadwick et al.,
 115 2003), and Keynote (Blaze et al., 1999) provide.
 116 The Globus team realised that ultimately this is
 117 what is needed for Grid authorisation, and so,
 118 within the remit of the Global Grid Forum (GGF),
 119 set about defining a standard interface between
 120 a Grid application that wants to know if a user is
 121 authorised to perform a certain action, and an
 122 authorisation infrastructure that is able to answer
 123 such questions.

124 An authorisation interface for the Grid

125 The ISO Access Control Framework standard (ITU-T
 126 Rec X.812, 1995) recognised nearly a decade ago
 127 that access control can be split into two compo-
 128 nents: an application dependent enforcement
 129 function (AEF) and an application independent
 130 decision function (ADF) – termed, respectively,
 131 the policy enforcement point (PEP) and policy
 132 decision point (PDP) in various IETF documents.
 133 The PDP (or ADF) can be controlled by a policy (or
 134 set of rules), and providing the policy is general
 135 enough, the PDP is able to make decisions for any
 136 type of application. All the ADF needs in addition
 137 to the policy is details about the user (initiator),
 138 the resource being protected (target), the access
 139 request and environmental (or contextual) infor-
 140 mation such as the time of day (see Fig. 1). In
 141 addition, if the ADF/PDP retains information about
 142 previous requested actions it can make subsequent
 143 decision based on this, for example, ATM machines
 144 that will only allow you to withdraw a maximum
 145 amount of money each day and will refuse re-
 146 quests once this limit has been reached. The Open

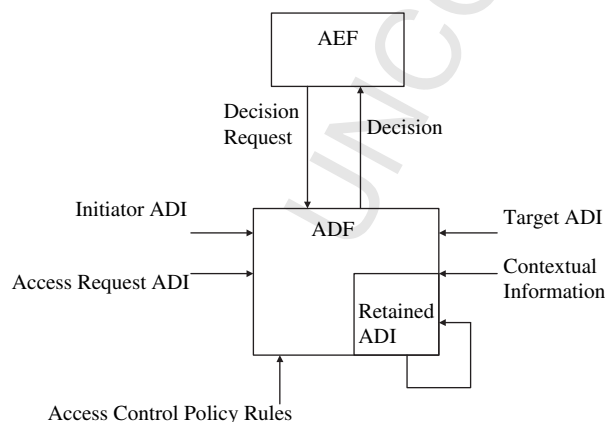


Figure 1 Separating access control enforcement from decision making.

147 Group defined an Application Programmable In-
 148 terface between the PEP and PDP and called it the
 149 AZN API (The Open Group, 2000). However, this is
 150 not general purpose enough, being constrained
 151 to C programs only.

152 The Grid is increasingly moving towards web
 153 services as the means of connecting its various
 154 components together. Globus Toolkit is moving this
 155 way too, and release 4 scheduled for the end of
 156 2004, is expected to be web services compliant. In
 157 2003 it was therefore opportune to specify a web
 158 services interface that could be used to connect
 159 a PDP to a Grid based PEP such as Globus Toolkit.
 160 Fortunately the groundwork for this had already
 161 been done by the Organization for the Advance-
 162 ment of Structured Information Standards (OASIS),
 163 who had issued the Security Assertion Markup
 164 Language (SAML) specification in 2002 (OASIS,
 165 2002a). SAML is a general purpose language that
 166 allows different types of security assertions about
 167 principals to be passed between clients and serv-
 168 ers, encoded as XML messages. The language is
 169 infinitely extensible and allows any type of asser-
 170 tion to be defined, although three standard types
 171 of assertions about principals were specified in
 172 SAML v1.0: authentication assertions, attribute
 173 assertions and authorisation decision assertions.
 174 It is the latter type that are passed between the
 175 PDP and PEP. SAML therefore provides a solid base
 176 from which to specify the PDP–PEP interface for
 177 Grid applications. The SAML specification also
 178 defines a request–response protocol in SOAP over
 179 HTTP for carrying the SAML assertions. SAML is thus
 180 fully web services compliant.

181 Whilst SAML provides a solid basis for specifying
 182 the PDP–PEP interface, it does not define every-
 183 thing that is needed for Grid applications. A SAML
 184 profile is thus needed to rectify the deficiencies,
 185 and this has been specified by the Global Grid
 186 Forum as a draft standard specification (Von Welch
 187 et al., 2004). Several important restrictions or
 188 additions to SAML have been specified, including:

- The contents of the authorisation response. A
 59 simple Boolean (granted or denied) is sufficient
 60 for the PDP–PEP interface, but SAML does not
 61 contain such a response,¹ hence the GGF
 62 profile defines one.
 63
- How the PDP gains access to the user’s author-
 64 isation credentials (initiator ADI in Fig. 1). Two
 65
 66

¹ The SAML Authorization Decision Response repeats the entire contents of the Authorisation Decision Request, which is useful if the request is sent by a party other than the PEP, for example, the principal. The SAML response thus details exactly what has been granted to the principal.

67 modes are possible, pull and push. In pull mode
 68 the PDP fetches the credentials, in push mode
 69 the PEP provides them. If the PDP is to fetch
 70 the credentials, how does it know where to get
 71 them from? The PDP could either be pre-
 72 configured with a (probably static) list of
 73 credential sources, or the client could tell the
 74 server where to pull the credentials from at the
 75 time of decision making. The latter is more
 76 scalable, and more dynamic than the former.
 77 Thus the GGF profile defines a Reference
 78 Statement which points to a repository where
 79 user credentials are located.

80 - Default values for all the parameters of the
 81 authorisation decision request. These are
 82 specified in the GGF draft as follows: if the
 83 name of the initiator is missing it is assumed to
 84 be anyone i.e. public access is being re-
 85 quested; if the requested action is missing it
 86 is assumed to be everything i.e. all the rights
 87 that have been granted to the initiator; if the
 88 target is missing it is assumed to be all the
 89 resources that are protected by the PDP policy;
 90 if the initiator's credentials are missing then
 91 only the default ones (if any) that have been
 92 granted to everyone should be used. Note that
 93 no default values for the contextual informa-
 94 tion have been specified since in general it is
 95 not possible to define these.

96 - If too little information is passed to the PDP, it
 97 may simply deny access. Alternatively, at its
 98 discretion, the PDP may return "Granted sub-
 99 ject to" along with a set of conditions that
 100 must be fulfilled before access is granted, e.g.,
 101 Granted subject to the time being between 9
 102 am and 5 pm. It is then the responsibility of the
 103 PEP to evaluate these conditions before grant-
 104 ing access to the initiator. If the PEP is unable
 105 or unwilling to evaluate these conditions, it
 106 always has the option of issuing a new decision
 107 request and sending more information to
 108 the PDP (such as the missing contextual
 109 information).

239 Once this interface had been defined by the
 240 GGF, it then needed to be implemented and tested
 241 in one or more Grid applications to ensure that it
 242 meets the needs of the Grid community. Globus
 243 Toolkit v3.3, released in April 2004, has imple-
 244 mented this SAML interface, as has the PERMIS
 245 authorization infrastructure, described in the next
 246 section. The BRIDGES E-Science project currently
 247 running at Glasgow University is the first Grid
 248 application to pilot the combined GT3.3/PERMIS
 249 infrastructure and the results are expected to be
 250 published the last quarter of 2004.

The PERMIS authorisation infrastructure 251

252 PERMIS is software developed under the EC PERMIS
 253 project (www.permis.org). It is now part of the US
 254 National Science Foundation's Middleware Initia-
 255 tive (NMI) software release (www.nsf-middleware.org). PERMIS is an attribute based access control
 256 (ABAC) infrastructure. ABAC is a superset of role
 257 based access controls (RBAC), in which access
 258 control decisions are made based upon any attrib-
 259 utes held by the user, and not just upon their
 260 organisational roles (as in conventional RBAC). In
 261 PERMIS, user attributes are held in X.509 standard
 262 attribute certificates (ACs) ([ISO 9594-8/ITU Rec.
 263 X.509, 2001](http://www.iso.org/iso/9594-8/ITU_Rec_X.509_2001)). An attribute certificate is a data
 264 structure that binds details about the holder to the
 265 attributes that are assigned to them, digitally
 266 signed by the issuing attribute authority. The AC
 267 is therefore tamper-proof, and its validity can be
 268 checked by validating its digital signature, and
 269 checking that it has not been previously revoked in
 270 the current revocation list.
 271

272 In PERMIS, managers throughout a VO can act as
 273 attribute authorities and assign attributes (in the
 274 form of X.509 ACs) to their staff – they do not
 275 need any prior permission to do this. All they need
 276 is a private signing key and a corresponding X.509
 277 public key certificate (plus the necessary software
 278 to create ACs). Thus the allocation of entitlements
 279 (or ACs) is distributed throughout the entire VO
 280 (and beyond if necessary²). However, each re-
 281 source owner, when setting the policy that con-
 282 trols access to his/her Grid resource, states which
 283 attribute authorities (s)he trusts to issue X.509
 284 ACs, and the PERMIS PDP will then discard all ACs
 285 presented to it that are not digitally signed by one
 286 of these trusted authorities. Thus we have partly
 287 accomplished one of our earlier stated goals, i.e.
 288 that a resource owner should be able to specify
 289 a policy for controlling access to his/her resource,
 290 and then leave the authorisation infrastructure to
 291 enforce it.

292 The PERMIS distribution contains three software
 293 tools for creating X.509 ACs. A user friendly graph-
 294 ical Attribute Certificate Manager (see [Fig. 2](#))
 295 is designed to make it very easy for managers to
 296 assign basic ACs to their staff, one by one. A more
 297 sophisticated Privilege Allocator can create more
 298 complex ACs, whilst a bulk loader tool is designed to
 299 allow large numbers of users to be automatically

² PERMIS does not restrict who can issue ACs. Thus, for example, a professional society such as the Law Society or the General Medical Council, could issue "lawyer" or "doctor" ACs to their members, and access control decisions could subsequently be based upon them.

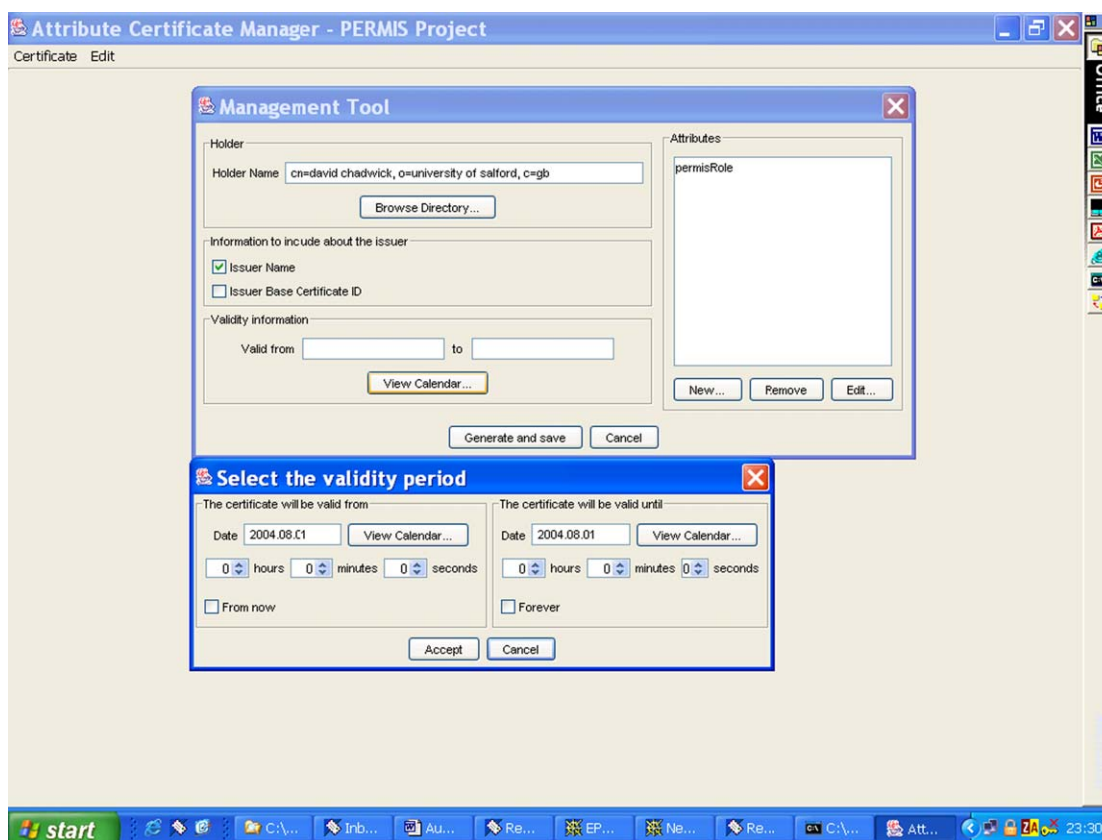


Figure 2 The Attribute Certificate Manager Tool.

300 allocated ACs. The bulk loader works by searching
 301 an LDAP directory for users who have specific
 302 attributes, then creating a specific AC for each of
 303 them and writing these back to their respective
 304 LDAP entries.

305 PERMIS provides a PDP that reads in the policy
 306 set by the resource owner at initialisation time. It
 307 then makes access control decisions for each
 308 authorisation decision request provided by the
 309 PEP, based on this policy. The interface between
 310 the PDP and the PEP is implemented as a Java API
 311 (for applications that want to combine the PDP and
 312 PEP as one program), and as SAML requests over
 313 SOAP and HTTP (for stand alone PDP servers, as
 314 used by GT3.3).

315 The user's X.509 ACs can be pushed to the
 316 PERMIS PDP by the PEP, either as complete X.509
 317 ACs, or as SAML Reference Statements. The PDP
 318 can also pull X.509 ACs from a set of pre-configured
 319 LDAP servers, or from the locations specified
 320 in the Reference Statements.

321 PERMIS policies are written in XML, according to
 322 a DTD published at www.xml.org. This DTD pre-
 323 dates the XACML specification (OASIS, 2000b), and
 324 for the most part is a subset of XACML, except that
 325 the PERMIS policy supports delegation of authority,
 326 a feature not currently supported by XACML. The

PERMIS authorisation policy is a set of sub-policies,
 327 namely:

- SubjectPolicy – this specifies the valid subject
 110 domains i.e. only users from these subject
 111 domains may be authorised to access resources
 112 covered by this policy. Each domain is specified
 113 as an LDAP subtree, using Include DN and
 114 Exclude DN statements. In Grid environments
 324 each user has a unique DN contained in his
 325 public key certificate. If his DN is not within
 326 a valid subject domain, the user will be denied
 327 all access to resources in the VO covered by
 328 this policy.
 329
- RoleHierarchyPolicy – this specifies the differ-
 330 ent roles and attributes that can be allocated
 331 to users, and the hierarchical relationships (if
 332 any) between them. A superior role inherits all
 333 the privileges of a subordinate role, as in
 334 conventional RBAC. Using role hierarchies can
 335 simplify Role Assignment Policies.
 336
- SOAPolicy – this specifies which attribute
 337 authorities (or Sources Of Authority) are
 338 trusted to allocate roles and attributes to
 339 users. This is the way that a resource owner
 340 specifies who within (and without) the VO is to
 341 be trusted to issue ACs.
 342

- 343 • RoleAssignmentPolicy – this specifies which
 344 roles and attributes may be allocated to which
 345 subjects by which SOAs, whether delegation of
 346 authority may take place or not, and how long
 347 the issued ACs are considered valid by this
 348 policy. This sub-policy effectively states who is
 349 trusted to allocate which roles to whom, and is
 350 central to the distributed management of
 351 trust. It can stop a trusted manager in one
 352 organisation issuing attributes or roles to staff
 353 in another organisation. Allowing delegation
 354 will allow a staff member to assign his/her
 355 attributes to another staff member. By re-
 356 stricting the validity of ACs, an issued AC may
 357 have a validity period of 2 years, but the
 358 resource owner may have a more stringent
 359 policy and be not prepared to accept ACs older
 360 than 1 year.
- 361 • TargetPolicy – this specifies the target domains
 362 in which resources covered by this policy are
 363 located. Each domain is specified as either an
 LDAP subtree, using Include DN and Exclude DN
 statements, or as a URL. Resources can further
 be refined by specifying their type e.g. all
 fileservers within the o = Salford, c = GB
 domain.
- 369 • ActionPolicy – this specifies the actions (or
 370 methods) supported by the various target
 371 resources, along with the parameters that
 372 should be passed along with each request e.g.
 373 action Open with parameter Filename.
- 374 • TargetAccessPolicy – this specifies which roles
 375 and attributes are needed in order to perform
 376 which actions on which targets, and under
 377 which conditions. Note that this part of the
 378 policy, being ABAC, is not concerned with the
 379 distinguished names of the users. Conditions
 380 are specified using Boolean logic and might
 381 contain constraints such as "IF time is GT 9 am
 382 AND time is LT 5 pm OR IF Calling IP address is
 383 a subset of 125.67.x.x". All actions that are not
 384 specified in a Target Access Policy are denied.

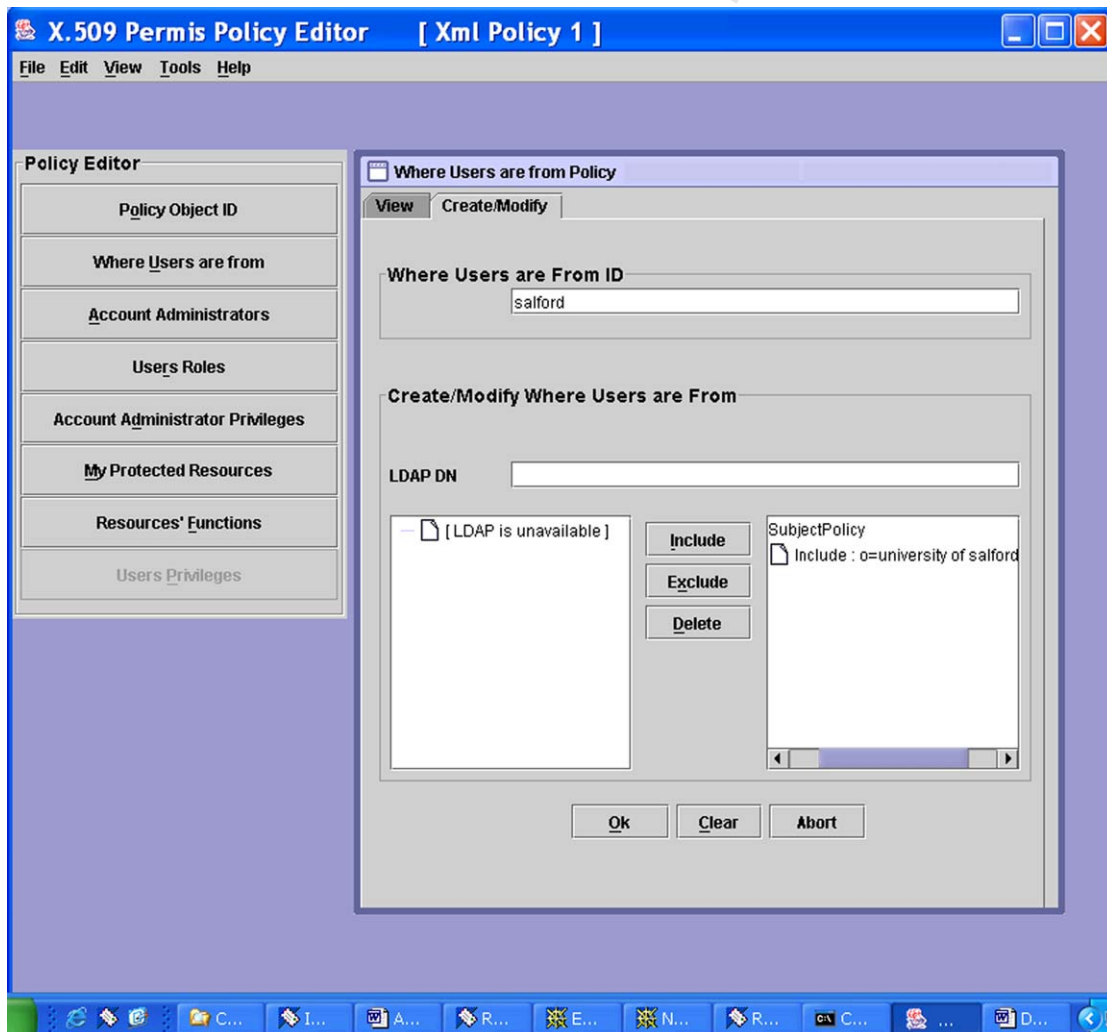


Figure 3 The Policy Management Tool.

395 Clearly setting the policy for controlling access
 396 to a resource is a potentially complex task,
 397 especially if the resource owner was required to
 398 write the XML to specify it. This would be beyond
 399 the capabilities of most resource owners. Conse-
 400 quently, a user friendly GUI is provided to make it
 401 easy to write PERMIS policies (see Fig. 3). Once the
 402 policy is written, the resource owner digitally signs
 403 it (to prevent it from being subsequently tampered
 404 with), and the GUI stores it in the owner's LDAP
 405 directory entry. Each policy also has a unique ID so
 406 that an owner can have several different policies
 407 for different occasions. When the PERMIS PDP is
 408 started it is told which policy ID to use, so it reads
 409 in this policy from the owner's LDAP entry, checks
 410 the signature and validity time, and if valid, the
 411 PDP knows that it has the correct policy written by
 412 the resource owner. The resource owner can then
 413 leave the authorisation infrastructure to control
 414 access to his/her resource without having to
 415 continually manage the system.

416 Conclusions and future work

417 Allowing a resource owner to set the policy for who
 418 has permission to perform which actions on which
 419 resource, and distributing the assignment of rights
 420 to managers throughout a VO, is clearly something
 421 that is needed to facilitate large scale Grid com-
 422 puting. Separating access control enforcement
 423 into a policy controlled application independent
 424 PDP and application dependent PEP, linked togeth-
 425 er via a web services SAML interface is something
 426 that will facilitate the development of compre-
 427 hensive and sophisticated PDPs such as PERMIS and
 428 Akenti. As Grid resource owners demand more
 429 access control features, such as separation of
 430 duties, mutually exclusive roles, dynamic delega-
 431 tion of authority, etc., then it becomes a question
 432 of adding suitable rules to the policy base and
 433 enhancing the PDP to correctly evaluate them.

434 At this point in time it is not clear which XML
 435 based policy language will become the de-facto
 436 standard for PDPs. XACML (OASIS, 2000b), from
 437 OASIS has some industry support, and Grid re-
 438 searchers are now experimenting with it. But
 439 Microsoft is also championing XrML (www.xrml.org)
 440 as a general purpose language for expressing who
 441 has the right to do what with digital content.
 442 Clearly there is significant overlap between these
 443 two languages, but with XrML now being part of
 444 Microsoft Office 2003 and Windows Server 2003, it
 445 could very well become the de-facto standard
 446 for authorisation. The downside of XrML is that

a number of its features are patented, so users will
 need to pay a license fee to use it.

Future work inside the GGF is likely to specify
 a management interface to the PDP that will allow
 the resource manager to dynamically update the
 policy that is to be used for decision making. In
 current systems it is a local matter how the PDP is
 configured with the correct policy to be used, and
 how the policy is changed according to the chang-
 ing circumstances in the VO. Further, in the future,
 autonomic PDPs could well be developed and
 configured with meta-policies telling them which
 access control policy to use under which condi-
 tions. As one can see, policy based authorisation
 and access controls are definitely here to stay.

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- Professor David Chadwick is the leader of the Information
 Systems Security Research Centre (ISSRC) at the University of

504 Salford. This is part of Informatics Research Institute (see 510
505 <http://www.iris.salford.ac.uk/>) which gained a 5* rating in the 511
506 December 2001 RAE. Prof Chadwick has written over 40 books, 512
507 journal and conference papers and the latest of these can be 513
508 downloaded from <http://sec.isi.salford.ac.uk/Papers>. He was 514
the editor of X.518 (1993) and is the author of several current 515
Internet and Global Grid Forum Draft Standards including the
Grid authorisation API draft standard. He has been the principal
investigator in over 10 research grants from a variety of sources
including JISC, EPSRC and the EC. He was the technical director
of the EC PERMIS project (www.permis.org) and instrumental in
its design and evolution.

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