Computer Science at Kent

Algorithmic Debugging for Locally Defined Functions

Yong Luo and Olaf Chitil

Technical Report No. 8 - 07
August 2007

Copyright © 2007 University of Kent
Published by the Computing Laboratory,
University of Kent, Canterbury, Kent, CT2 7NF, UK
Algorithmic Debugging for Locally Defined Functions

Yong Luo and Olaf Chitil
Computing Laboratory, University of Kent

Abstract The purpose of the document is to prove the correctness of Algorithmic Debugging where the traces for local functions are generated in a new way. The processes of generating computation graphs follow exactly what we might do by hand. Therefore, we can be confident that the graphs are correct. We do not need to justify the graphs by comparing $\lambda$-lifted programs.

1 Basic Definitions

In this section we give some basic definitions.

Definition 1. (Nodes, Atoms)

- A node is a sequence of letters r, f and a, i.e. \{r,f,a\}*. 
- Atoms:
  1. a constructor is an atom; 
  2. a function symbol is an atom; 
  3. a node combined with a function symbol is an atom. For example, $m.f$ is an atom where $m$ is a node and $f$ is a function symbol.

Notation: In the future, we shall say that $g$ is a function if $g$ is an atom but not a constructor.

Definition 2. (Terms, Patterns, Rewriting rule and Program)

- Terms:
  1. an atom is a term; 
  2. a node is a term; 
  3. a variable is a term; 
  4. (Application) $MN$ is a term if $M$ and $N$ are terms.

- Patterns:
  1. a variable is a pattern; 
  2. $cp_1...p_n$ is a pattern if $c$ is a constructor and $p_1,...,p_n$ are patterns, and the arity of $c$ is $n$. 

A **simple rewriting rule** is of the form \( f \, p_1\ldots p_n = R \) where \( f \) is a function and \( p_1, \ldots, p_n (n \geq 0) \) are patterns and \( R \) is a term.

A **rewriting rule** is in one of two forms:
1. (top-level functions without local functions) a simple rewriting rule.
2. (top-level functions with local functions) the form

\[
    f \, p_1\ldots p_n = R \\
    \text{where } g_1 \, q_{11}\ldots q_{m_1} = R_1 \\
    \text{ ..... } \\
    g_k \, q_{1k}\ldots q_{m_k} = R_k
\]

where \( f \, p_1\ldots p_n = R \) and \( g_j \, q_{1j}\ldots q_{m_j} = R_j \) are simple rewriting rules, \( g_1, \ldots, g_k \) are the local functions of \( f \).

A **program** is a finite set of rewriting rules.

If a simple rewriting rule is of the form \( f = R \) we call it a constant rewriting rule and \( f \) is a constant.

**Definition 3. (Node expression and Computation graph)**

- A **node expression** is either
  - an atom, or
  - a node, or
  - an application of two nodes, which is of the form \( m \circ n \).

- A **computation graph** is a set of pairs which are of the form \( (n, e) \), where \( n \) is a node and \( e \) is a node expression.

**Notation:** \( \text{dom}(G) \) denotes the set of nodes in a computation graph \( G \).

### 2 Pattern matching

The pattern matching algorithm for a graph has two different results, either a set of substitutions or “doesn’t match”.

- Let \( G \) be a computation graph, and \( m \in \text{dom}(G) \). The final node in a sequence of reductions starting at \( m \), \( \text{last}_G(m) \):

\[
    \text{last}_G(m) = \begin{cases} 
    \text{last}_G(mr) & \text{if } mr \in \text{dom}(G) \\ 
    \text{last}_G(n) & \text{if } (m, n) \in G \text{ and } n \text{ is a node} \\ 
    m & \text{otherwise}
    \end{cases}
\]

The purpose of this function is to find out the most evaluated point for \( m \).
Let $G$ be a computation graph, and $m \in \text{dom}(G)$. The head of the term at $m$, $\text{head}_G(m)$:

$$
\text{head}_G(m) = \begin{cases} 
\text{head}_G(\text{last}_G(i)) & \text{if } (m, i \circ j) \in G \\
a & \text{if } (m, a) \in G \text{ and } a \text{ is an atom} \\
\text{undefined} & \text{otherwise}
\end{cases}
$$

Let $G$ be a computation graph, and $m \in \text{dom}(G)$. The arguments of the function at $m$, $\text{args}_G(m)$, is defined as follows.

$$
\text{args}_G(m) = \begin{cases} 
\langle \text{args}_G(\text{last}(G, i)), j \rangle & \text{if } (m, i \circ j) \in G \\
\emptyset & \text{otherwise}
\end{cases}
$$

Note that the arguments of a function are a sequence of nodes.

Now, we define two functions $\text{match}_1$ and $\text{match}_2$ which are mutually recursive. The arguments of $\text{match}_1$ are a node and a pattern. The arguments of $\text{match}_2$ are a sequence of nodes and a sequence of patterns.

- $\text{match}_1$:
  
  $$
  \text{match}_1_G(m, x) = [m/x] \text{ where } x \text{ is a variable}
  
  \text{match}_1_G(m, q_1...q_k)
  = \begin{cases} 
  \text{match}_2_G(\text{args}(G, m'), \langle q_1, ..., q_k \rangle) & \text{if } \text{head}_G(m') = c \\
  \text{does not match} & \text{otherwise}
  \end{cases}
  $$

  where $m' = \text{last}_G(m)$.

- $\text{match}_2$:
  
  $$
  \text{match}_2_G(\langle m_1, ..., m_n \rangle, \langle p_1, ..., p_n \rangle)
  = \text{match}_1_G(m_1, p_1) \cup ... \cup \text{match}_1_G(m_n, p_n)
  $$

  where $\cup$ is the union operator. Notice that if $n = 0$ then

  $$
  \text{match}_2_G(\emptyset, \emptyset) = \emptyset
  $$

  If any $m_i$ does not match $p_i$, $\langle m_1, ..., m_n \rangle$ does not match $\langle p_1, ..., p_n \rangle$. If the length of two sequences are not the same, they do not match. For example, $\langle m_1, ..., m_s \rangle$ does not match $\langle p_1, ..., p_{s'} \rangle$ if $s \neq s'$.

- We say that $G$ at $m$ matches the left-hand side $fp_1...p_n$ of a rewriting rule with $[m_1/x_1, ..., m_k/x_k]$ if $\text{head}_G(m) = f$ and

  $$
  \text{match}_2_G(\text{args}_G(m), \langle p_1, ..., p_n \rangle) = [m_1/x_1, ..., m_k/x_k]
  $$

  In the substitution form $[m/x]$, $m$ is not a term but a node. The definition of pattern matching and its result substitution sequence will become important for making computation order irrelevant when we generate graphs.
3 Renaming and Program for local functions

Suppose that $G$ at $m$ matches the left-hand side of a rewriting rule with $[m_1/x_1, ..., m_l/x_l]$, and the rewriting rule has local functions as follows.

\[
\begin{align*}
    f \ p_1...p_n &= R \\
    \text{where } g_1 \ q_1...q_{r_1} &= R_1 \\
    \vdots \\
    g_k \ q_{k_1}...q_{r_k} &= R_k
\end{align*}
\]

We generate a set of new simple rewriting rules, $L_m$, called local functions at $m$. All the local functions $g_1, ..., g_k$ in the local rewriting rules are renamed to $m.g_1, ..., m.g_k$, and all the free variable in $R_1, ..., R_k$ are substituted by $[m_1/x_1, ..., m_l/x_l]$. Then, $L_m$ looks like the following:

\[
\begin{align*}
    m.g_1 \ q_1...q_{r_1} &= R'_1 \\
    \vdots & \vdots \\
    m.g_k \ q_{k_1}...q_{r_k} &= R'_k
\end{align*}
\]

4 ART

The function $\text{graph}$ is defined as follows.

**Definition 4.** ($\text{graph}$) Let $G$ be a computation graph, and $m \in \text{dom}(G)$. The function $\text{graph}$ takes two arguments. The first argument is a node and the second is a term.

\[
\text{graph}(n, e) = \begin{cases} 
(n, e) & \text{where } e \text{ is an atom or a node} \\
\{(n, M \circ N)\} & \text{if } M \text{ and } N \text{ are nodes} \\
\{(n, M \circ n a)\} \cup \text{graph}(n a, N) & \text{if only } M \text{ is a node} \\
\{(n, n f \circ N)\} \cup \text{graph}(n f, M) & \text{if only } N \text{ is a node} \\
\{(n, n f \circ n a)\} \cup \text{graph}(n f, M) & \text{otherwise} \\
\cup \text{graph}(n a, N) & 
\end{cases}
\]

**Generate an ART**

- For a starting term $M$, the starting ART is $\text{graph}(r, M)$. Note that the start term has no nodes inside.
- (*ART rule 1*) If an ART $G$ at $m$ matches the left-hand side of a simple rewriting rule $f p_1...p_n = R$ in the program $L$ with $[m_1/x_1, ..., m_l/x_l]$, then we generate a new ART.

\[
G \cup \text{graph}(mr, R[m_1/x_1, ..., m_l/x_l])
\]
• **(ART rule 2)** If an ART \( G \) at \( m \) matches the left-hand side of a rewriting rule in the program \( L \) with \([m_1/x_1, \ldots, m_l/x_l]\), and the rewriting rule has local functions as follows.

\[
\begin{align*}
  f \ p_1 \ldots p_n &= R \\
  \text{where } g_1 q_1 \ldots q_{r_1} &= R_1 \\
  \ldots \ldots \\
  g_k q_1 \ldots q_{r_k} &= R_k
\end{align*}
\]

Then we generate a set of new rewriting rules \( L_m \) and a new ART.

\[
G \cup \text{graph}(mr, R'[m_1/x_1, \ldots, m_l/x_l])
\]

where \( R' \) is obtained from \( R \) by renaming all the local functions in \( R \).

• **(ART rule 3)** If an ART \( G \) at \( m \) matches the left-hand side of a simple rewriting rule \((s.f)p_1 \ldots p_n = R\) in the program \( L_s \) with \([m_1/x_1, \ldots, m_k/x_k]\), then we generate a new ART.

\[
G \cup \text{graph}(mr, R[m_1/x_1, \ldots, m_k/x_k])
\]

• An ART is generated from the starting ART and by applying the ART rules repeatedly. Note that the order in which nodes are chosen to compute has no influence in the final graph.

The following simple properties of an ART will be used later.

**Lemma 1.** Let \( G \) be an ART.

- If \( m \in \text{dom}(G) \) then there is at least one letter \( r \) in \( m \).
- If \( mr \in \text{dom}(G) \) then \( m \in \text{dom}(G) \) or \( m = \varepsilon \) where \( \varepsilon \) is the empty sequence.
- If \( mr \in \text{dom}(G) \) then \((m, n) \notin G \) for any node \( n \).

**Proof.** The first and second are trivial. The third is proved by contradiction. If \((m, n) \in G \) then \( \text{head}_G(m) \) is undefined. There cannot be a computation at \( m \), i.e. \( mr \notin G \).

5 EDT

Generating an Evaluation Dependency Tree

**Definition 5.** *(Parent edges)*

\[
\begin{align*}
\text{parent}(nf) &= \text{parent}(n) \\
\text{parent}(na) &= \text{parent}(n) \\
\text{parent}(nr) &= n
\end{align*}
\]
Note that $parent(r) = \varepsilon$ where $\varepsilon$ is the empty sequence.

**Definition 6. (children and tree)** Let $G$ be an ART, and $mr$ a node in $G$ (i.e. $mr \in \text{dom}(G)$). children and tree are defined as follows.

- **children:**
  
  $$children(m) = \{n \mid nr \in \text{dom}(G) \text{ and } parent_G(n) = m\}$$

- **tree:**
  
  $$tree(m) = \{(m, n_1), ..., (m, n_k)\} \cup tree(n_1) \cup ... \cup tree(n_k)$$

  where $\{n_1, ..., n_k\} = children(m)$

Usually, a single node of a computation graph represents many different terms. We are particularly interested in two kinds of terms of nodes, the most evaluated form and the redex.

**Definition 7. (Most Evaluated Form)** Let $G$ be an ART. The most evaluated form of a node $m$ is a term and is defined as follows.

$$mef(m) = \begin{cases} 
  \text{mef}(mr) & \text{if } mr \in \text{dom}(G) \\
  \text{mef}(m) & \text{otherwise}
\end{cases}$$

where

$$mef(m) = \begin{cases} 
  a & (m, a) \in G \text{ and } a \text{ is an atom} \\
  \text{mef}(n) & (m, n) \in G \text{ and } n \text{ is a node} \\
  \text{mef}(i) \text{mef}(j) (m, i \circ j) \in G
\end{cases}$$

One may also use the definition of $last_G(m)$ to define the most evaluated form.

**Definition 8. (redex)** Let $G$ be an ART, and $mr$ a node in $G$ (i.e. $mr \in \text{dom}(G)$). redex is defined as follows.

- **redex($\varepsilon$) = main**

- **redex($m$) =**
  
  $$\begin{cases} 
  \text{mef}(i) \text{mef}(j) & (m, i \circ j) \in G \\
  a & (m, a) \in G \text{ and } a \text{ is an atom}
\end{cases}$$
Generate an EDT

Now, we define the evaluation dependency tree of a graph.

**Definition 9. (Evaluation Dependency Tree)** Let $G$ be an ART. The evaluation dependency tree (EDT) of $G$ consists of the following two parts.

1. The set $\text{tree}(\varepsilon)$;
2. The set of equations; for any node $m$ in $\text{tree}(\varepsilon)$ there is a corresponding equation. If $\text{head}_G(m)$ is a top-level function, the the equation at $m$ is of the form
   \[ \text{redex}(m) = \text{mef}(m) \]
   If $\text{head}_G(m)$ is a local function of the form $n.f$, then the equation at $m$ is of the form
   \[ \text{redex}(m) = \text{mef}(m) \]
   within $\text{redex}(n)$

6 Proofs

Some of the definitions and proofs are as the same before but some are new. The old proofs still need to be checked again because the basic definitions such as rewriting rules and EDT are changed.

The following theorems suggest that the EDT of an ART covers all the computation in the ART. Although two evaluations may rely on the same evaluation in an ART, every evaluation for algorithmic debugging only needs to be examined once.

**Lemma 2.** Let $G$ be an ART, and $T$ its EDT. If there is a sequence of nodes $m_1, m_2, ..., m_k$ such that
\[
m \in \text{children}(m_1), m_1 \in \text{children}(m_2), ..., \\
   m_{k-1} \in \text{children}(m_k), m_k \in \text{children}(\varepsilon)
\]
then $m \in \text{dom}(T)$.

**Proof.** By the definition of $\text{tree}(\varepsilon)$.

**Lemma 3.** Let $G$ be an ART. If $m \in \text{dom}(G)$, then $m \equiv \varepsilon$ or there is a sequence of nodes $m_1, m_2, ..., m_k$ such that
\[
m \in \text{children}(m_1), m_1 \in \text{children}(m_2), ..., \\
   m_{k-1} \in \text{children}(m_k), m_k \in \text{children}(\varepsilon)
\]
Proof. By induction on the size of $m$, and by Lemma 1.

Since $mr \in \text{dom}(G)$, by Lemma 1, we only need to consider the following two cases.

- If $m = \varepsilon$, the statement is obviously true.
- If $m \in \text{dom}(G)$, by Lemma 1, there is at least one letter $r$ in $m$. We consider the following two sub-cases.
  
  · $m = rn$, where there is no $r$ in $n$. Since $mr \in \text{dom}(G)$ and $\text{parent}(rn) = \varepsilon$, we have $rn \in \text{children}(\varepsilon)$.
  · $m \equiv m_1rn$, where there is no $r$ in $n$. Since $mr \in \text{dom}(G)$ and $\text{parent}(m) = m_1$, we have $m \in \text{children}(m_1)$. Now, because $m_1$ is a sub-sequence of $m$, by induction hypothesis, there is a sequence of index numbers $m_2, \ldots, m_k$ such that

\[
m_1 \in \text{children}(m_2), \ldots, m_{k-1} \in \text{children}(m_k), m_k \in \text{children}(\varepsilon)
\]

So, there is a sequence of index numbers $m_1, m_2, \ldots, m_k$ such that

\[
m \in \text{children}(m_1), m_1 \in \text{children}(m_2), \ldots, m_k \in \text{children}(\varepsilon)
\]

**Theorem 1.** Let $G$ be an ART, and $T$ its EDT.
If $mr \in \text{dom}(G)$, then $m \in \text{dom}(T)$. In other words, $T$ covers all the computations in $G$.

*Proof.* By Lemma 3 and 2.

**Lemma 4.** Let $G$ be an ART, and $T$ its EDT.
If $(m, n) \in T$, then $n \in \text{children}(m)$ and $\text{parent}(n) \equiv m$.

*Proof.* By the definition of tree.

**Theorem 2.** Let $G$ be an ART, and $T$ its EDT.
If $(m, n) \in T$ and $m \neq k$, then $(k, n) \notin T$.

*Proof.* By Lemma 4.

The above theorem suggests that every evaluation for algorithmic debugging only needs to be examined once although two evaluations may rely on the same evaluation. For example, $g$ is defined as $g x = (\text{not } x, \text{not } x, \text{not } x)$. When we compute $g (\text{not True})$, the equation $\text{not True} = \text{False}$ only appears once in the EDT.
6.1 Semantical Equality

**Notations:** \( M \simeq_I N \) means that \( M \) is equal to \( N \) with respect to the semantics of the programmer’s intention. If the evaluation \( M = N \) of a node in an EDT is in the programmer’s intended semantics, then \( M \simeq_I N \). Otherwise, \( M \not\simeq_I N \) i.e. the node is erroneous.

**Remark 1.** For a local function, the evaluation of a node in an EDT is of the form
\[
(m.g)b_1, ..., b_n = N
\]
within \( f_{e_1, ..., e_k} \).

The “within” part helps the programmer to decide whether the evaluation is intended or not, but it will not be used in proofs. We keep the prefix \( m \) in order to make semantics of local functions clear. If both \( m \) and “within” are removed, we might have \( gb_1, ..., b_n = N \) and \( gb_1, ..., b_n = N' \) where \( N \) and \( N' \) are different. In practice, one may choose different ways to help the programmer to answer such questions.

6.2 Equivalent rewriting rules

Two kind of rewriting rules, top-level and local-level, are used during the processes of building trace. For a top level rewriting rule, there is no node in the right-hand side. However, for a local rewriting rule of the form \((m.f)p_1...p_n = R\), it is possible that there are nodes in \( R \). When the computation stops and we start to analysis the properties, we regard that the rewriting rule is equivalent to \((m.f)p_1...p_n = R'\) where \( R' \) is obtained from \( R \) by replacing all the nodes by their most evaluated forms. For example, the nodes in \( R \) are \( m_1, ..., m_k \), then
\[
R' \equiv R[mef(m_1)/m_1, ..., mef(m_k)/m_k]
\]

For a top-level rewriting rule \( fp_1...p_n = R \), if it used at node \( m \) and there are local functions in \( R \), we regard that the rewriting rule is equivalent to \( fp_1...p_n = R' \) where \( R' \) is obtained from \( R \) by renaming all the local functions in \( R \).

6.3 Program faulty

**Definition 10. (Program faulty)**

- For a simple rewriting rule \( fp_1...p_n = R \) without local functions, if there exists a substitution \( \sigma \) such that \((fp_1...p_n)\sigma \not\simeq_I R\sigma\), then we say that the definition of the function \( f \) in the program is faulty.
For a rewriting rule without local functions of the following form

\[ f \ p_1 \ldots p_n = R \]

where \( g_1 \ q_{11} \ldots q_{m1} = R_1 \)

\[ \ldots \]

\( g_k \ q_{1k} \ldots q_{mk} = R_k \)

If there exists a substitution \( \sigma \) such that \( (fp_1 \ldots p_n)\sigma \not\simeq_I (R \text{ within } fp_1 \ldots p_n)\sigma \), then we say that the definition of the function \( f \) in the program is faulty. If there exists \( \sigma \) and \( \sigma' \) such that \( \{(g_i \ q_{1i} \ldots q_{mi})\sigma' \not\simeq_I R_i\sigma' \} \text{ within } (fp_1 \ldots p_n)\sigma \), then we say that the definition of the local function \( g_i \) within \( f \) is faulty.

### 6.4 Correctness of Algorithmic Debugging

The proofs are the same as before, but we should check them again because the definition of program faulty is changed.

**Definition 11.** If the following statement is true, then we say that algorithmic debugging is correct.

- If the equation of a faulty node is \( fb_1 \ldots b_n = M \), then the definition of the function \( f \) in the program is faulty.

For a faulty node \( m \), we have \( \text{redex}(m) \not\simeq_I \text{mef}(m) \). We shall find a term \( N \) and prove \( \text{redex}(m) \rightarrow_P N \equiv_I \text{mef}(m) \). In order to define \( N \), we need other definitions.

**Definition 12.** Let \( G \) be an ART and \( m \) a node in \( G \). \( \text{reduct}(m) \) is defined as follows.

\[
\text{reduct}(m) = \begin{cases} 
  a & \text{if } (m, a) \in G \text{ and } a \text{ is an atom} \\
  \text{mef}(n) & \text{if } (m, n) \in G \text{ and } n \text{ is a node} \\
  \text{reduct}(mf) \text{ reduct}(ma) & \text{if } (m, mf \circ ma) \in G \\
  \text{reduct}(mf) \text{ mef}(j) & \text{if } (m, mf \circ j) \in G \text{ and } j \neq ma \\
  \text{mef}(i) \text{ reduct}(ma) & \text{if } (m, i \circ ma) \in G \text{ and } i \neq mf \\
  \text{mef}(i) \text{ mef}(j) & \text{if } (m, i \circ j) \in G \text{ and } i \neq mf \text{ and } j \neq ma 
\end{cases}
\]

\( \text{reduct} \) represents the result of a single-step computation. And we shall prove \( \text{redex}(m) \rightarrow_P \text{reduct}(mr) \equiv_I \text{mef}(m) \) for a faulty node \( m \). Note that \( \text{mef}(m) = \text{mef}(mr) \) and so we want to prove \( \text{reduct}(mr) \equiv_I \text{mef}(mr) \).

In order to prove this, we prove a more general result \( \text{reduct}(m) \equiv_I \text{mef}(m) \) for all \( m \in \text{dom}(G) \) (see Lemma 6 for the conditions).

We define \( \text{branch} \) and the reduction principle \( \text{depth} \) in order to prove this general result.
Definition 13. (*branch and branch*'') We say that $n$ is a branch node of $m$, denoted as $\text{branch}(n, m)$, if one of the following holds.

- $\text{branch}(m, m)$;
- $\text{branch}(n_f, m)$ if $\text{branch}(n, m)$;
- $\text{branch}(n_a, m)$ if $\text{branch}(n, m)$.

Let $G$ be an ART.

$$\text{branch}'(m) = \{ n \mid n_r \in \text{dom}(G) \text{ and } \text{branch}(n, m) \}$$

Note that $\text{branch}'(m)$ is the set of all evaluated branch nodes of $m$.

Lemma 5. Let $G$ be an ART.

- If $n \in \text{branch}'(m_f)$ or $n \in \text{branch}'(m_a)$ then $n \in \text{branch}'(m)$.
- If $m_r \in \text{dom}(G)$ then $\text{children}(m) = \text{branch}'(m_r)$.

Proof. By the definitions of $\text{children}$ and $\text{branch}'$.

Definition 14. (*depth*) Let $m$ be a node in an ART $G$.

$$\text{depth}(m) = \begin{cases} 
1 + \max\{ \text{depth}(m_f), \text{depth}(m_a) \}, & \text{if } (m, m_f \circ m_a) \in G \\
1 + \text{depth}(m_f), & \text{if } (m, m_f \circ j) \in G \text{ and } j \neq m_a \\
1 + \text{depth}(m_a), & \text{if } (m, i \circ m_a) \in G \text{ and } i \neq m_f \\
1, & \text{if } (m, i \circ j) \in G \text{ and } i \neq m_f \text{ and } j \neq m_a \\
0, & \text{otherwise}
\end{cases}$$

Lemma 6. Let $G$ be an ART and $m$ a node in $G$. If $\text{redex}(n) \simeq_I \text{mef}(n)$ for all $n \in \text{branch}'(m)$, then $\text{reduct}(m) \simeq_I \text{mef}(m)$.

Proof. By induction on $\text{depth}(m)$.

When $\text{depth}(m) = 0$, we have $(m, e) \in G$ where $e$ is a node or an atom.

- If $e$ is a node, then $m_r \in G$ by Lemma 1. Then by the definitions of $\text{reduct}$ and $\text{mef}$, we have $\text{reduct}(m) = \text{mef}(e)$ and $\text{mef}(m) = \text{meft}(m) = \text{mef}(e)$.
- If $e$ is an atom, we have $\text{reduct}(m) = e$. Now, we consider the following two cases. If $m \in \text{branch}'(m)$, then we have $m_r \in \text{dom}(G)$ and $\text{mef}(m) \simeq_I \text{redex}(m) = e$. If $m \notin \text{branch}'(m)$, then we have $m_r \notin \text{dom}(G)$ and $\text{mef}(m) = \text{meft}(m) = e$.

For the step cases, we proceed as follows.
• If \( m \in \text{branch}'(m) \), then we have \( mr \in \text{dom}(G) \) and \( \text{redex}(m) \simeq_I \text{mef}(m) \). And we need to prove \( \text{redex}(m) \simeq_I \text{reduct}(m) \).

Let us consider only one case here. The other cases are similar. Suppose \((m, mf \circ j) \in G \) and \( j \neq ma \), then by the definitions we have

\[
\text{redex}(m) = \text{mef}(mf) \text{mef}(j)
\]

\[
\text{reduct}(m) = \text{reduct}(mf) \text{mef}(j)
\]

Since for any \( n \in \text{branch}'(mf) \), by Lemma 5, we have \( n \in \text{branch}'(m) \) and hence \( \text{redex}(n) \simeq_I \text{mef}(n) \). By the definition of depth, we also have \( \text{depth}(mf) < \text{depth}(m) \). Now, by induction hypothesis, we have \( \text{reduct}(mf) \simeq_I \text{mef}(mf) \). And hence we have \( \text{redex}(m) \simeq_I \text{reduct}(m) \).

• If \( m \notin \text{branch}'(m) \), then \( mr \notin \text{dom}(G) \).

Let us also consider only one case. The other cases are similar. Suppose \((m, mf \circ j) \in G \) and \( j \neq ma \), then by the definitions we have

\[
\text{mef}(m) = \text{mef}(mf) \text{mef}(j)
\]

\[
\text{reduct}(m) = \text{reduct}(mf) \text{mef}(j)
\]

The same arguments as above suffice.

**Corollary 1.** Let \( G \) be an ART and \( mr \) a node in \( G \) (i.e. \( mr \in \text{dom}(G) \)).

If \( \text{redex}(n) \simeq_I \text{mef}(n) \) for all \( n \in \text{children}(m) \), then \( \text{reduct}(mr) \simeq_I \text{mef}(m) \).

**Proof.** By Lemma 5 and 6.

The condition, \( \text{redex}(n) \simeq_I \text{mef}(n) \) for all \( n \in \text{children}(m) \), basically means that \( m \) does not have any erroneous child nodes.

**Lemma 7.** Let \( G \) be an ART and \( mr \) a node in \( G \) (i.e. \( mr \in \text{dom}(G) \)).

Then \( \text{redex}(m) \rightarrow_P \text{reduct}(mr) \).

**Proof.** Since there is a computation at the node \( m \), we suppose \( G \) at node \( m \) matches the left-hand side of the rewriting rule \( fp_1...p_n = R \) with \( [m_1/x_1,...,m_k/x_k] \). We need to prove that there exists a substitution \( \sigma \) such that \( \text{redex}(m) = (fp_1...p_n)\sigma \) and \( \text{reduct}(mr) = R\sigma \). In fact \( \sigma = [\text{mef}(m_1)/x_1,...,\text{mef}(m_k)/x_k] \).

Now, we need to prove that \( \text{redex}(m) = (fp_1...p_n)\sigma \) and \( \text{reduct}(mr) = R\sigma \). For the first, we proceed by the definition of \( \text{redex} \) and pattern matching. For the second, we proceed by the definition of \( \text{reduct} \) and \( \text{graph} \).

Now, we come to the most important theorem, the correctness of algorithmic debugging.
Theorem 3. (Correctness of Algorithmic Debugging) Let $G$ be an ART, $T$ its EDT and $m$ a faulty node in $T$. If the equation for the faulty node $m$ is $fb_1...b_n = M$, then the definition of $f$ in the program is faulty.

Proof. By Lemma 7 and Corollary 1, we have $\text{redex}(m) \rightarrow_P \text{reduct}(mr)$ and $\text{reduct}(mr) \simeq_I \text{mef}(m)$. Since $fb_1...b_n \equiv \text{redex}(m) \not\simeq_I \text{mef}(m) \equiv M$, we have $fb_1...b_n \rightarrow_P \text{reduct}(mr)$ and $fb_1...b_n \not\simeq_I \text{reduct}(mr)$. The computation from $fb_1...b_n$ to $\text{reduct}(mr)$ is a single step computation, but $fb_1...b_n$ is not semantically equal to $\text{reduct}(mr)$. So the definition of $f$ in the program must be faulty.