

Transforming Taskell for Tracing

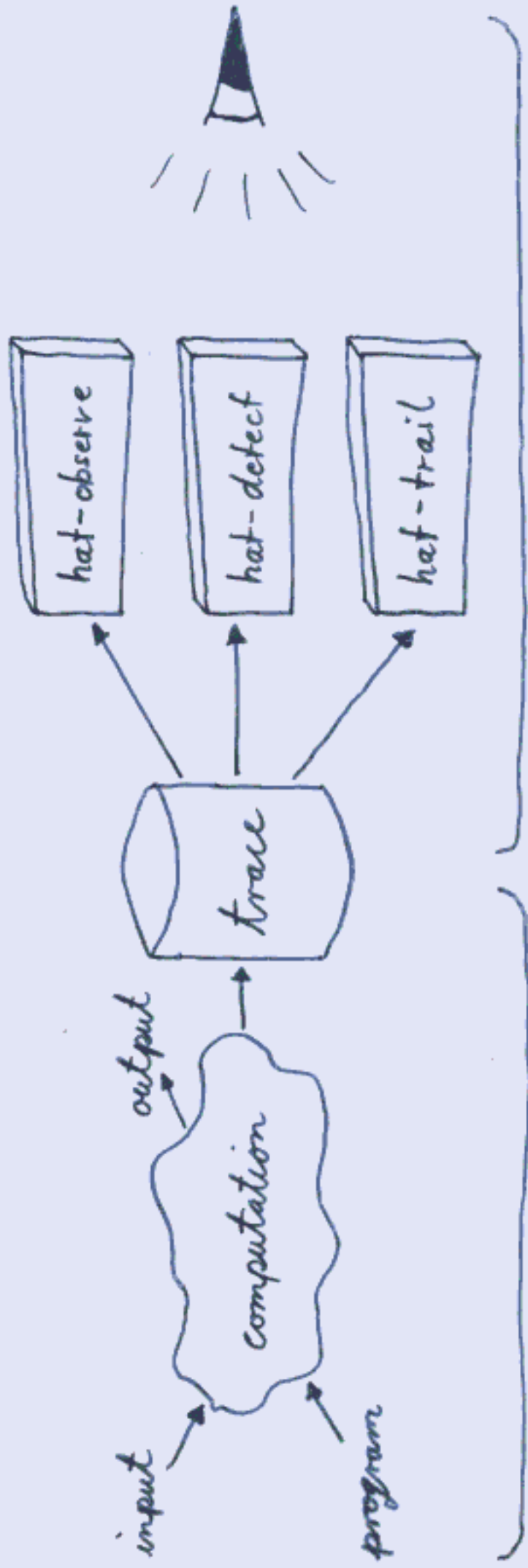
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Tracing with Hat

applications: program comprehension and debugging



trace as data structure liberates from the time arrow of the computation

Trace Generation through Transformation

was: Haskell $\rightarrow$ nhc98 $\rightarrow$ self-tracing executable

now: Haskell $\rightarrow$ hat-trans $\rightarrow$ Haskell source $\rightarrow$ self-tracing executable

Haskell compiler $\leftarrow$ Flat library

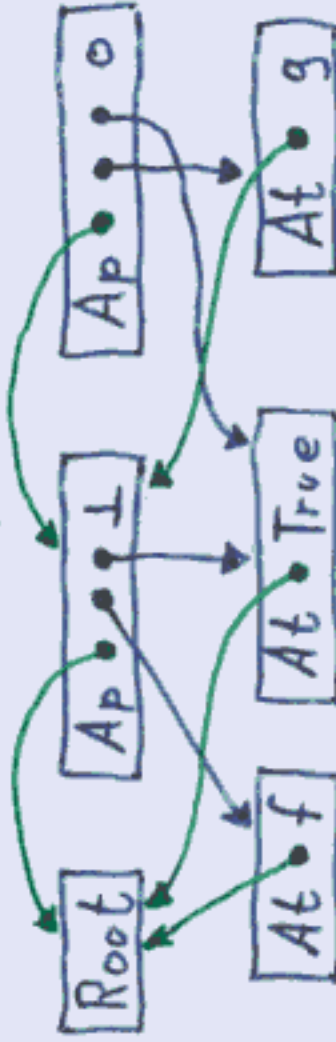
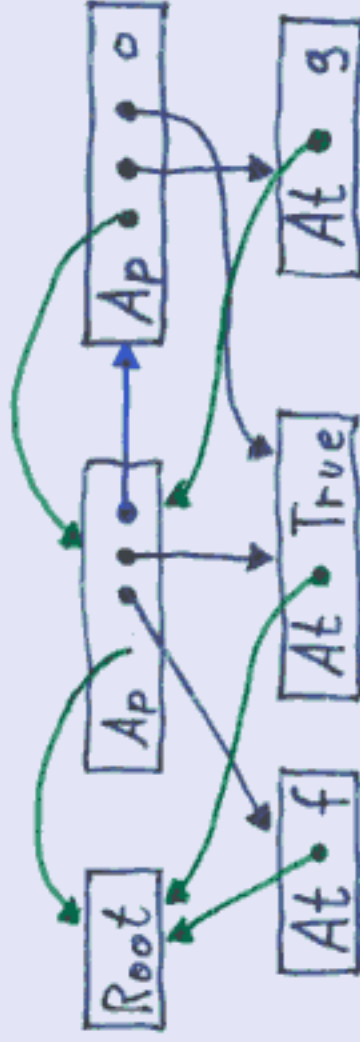
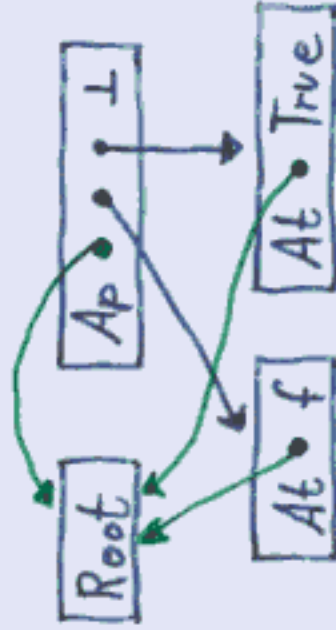
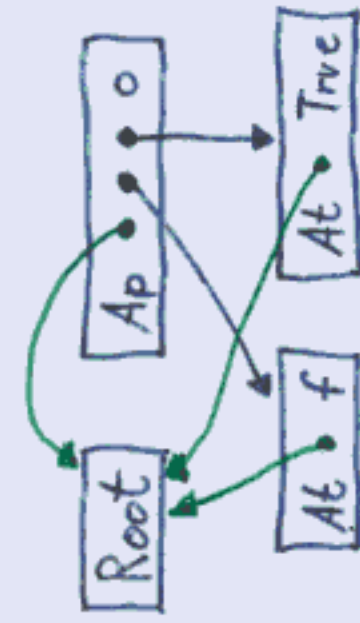
Why?

- comprehensible
- simplifies development
- combinable with different compilers

What does the trace of a reduction look like?

$$f \times = g \times$$

$$f \text{ True} \Rightarrow g \text{ True}$$

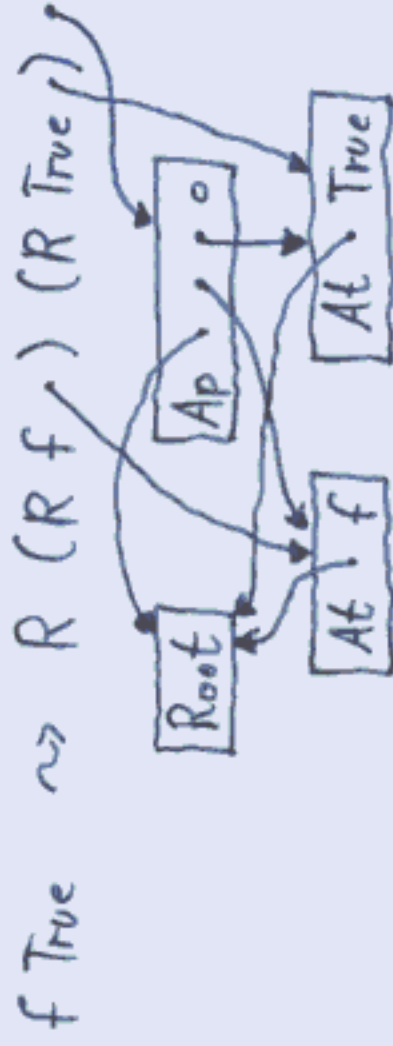


Do we use the IO-monad for writing the trace?

Do we use reflection to obtain meta-information such as identifier names?

How can evaluation of original expression and writing of trace be kept in synchrony?

data $R\ a = R\ a\ RefExp$



Does the transformation need type information?

Rat Library

- primitive combinators write trace file: mkAt, mkAp, entRedex, updResult, ...
- high-level combinators simplify the transformation: ap, ...

N-ary application and abstraction:

$$\text{ap1} :: \text{RefSrcPos} \rightarrow \text{RefExp} \rightarrow R (\text{Fun } a \ z) \rightarrow R a \rightarrow R z$$

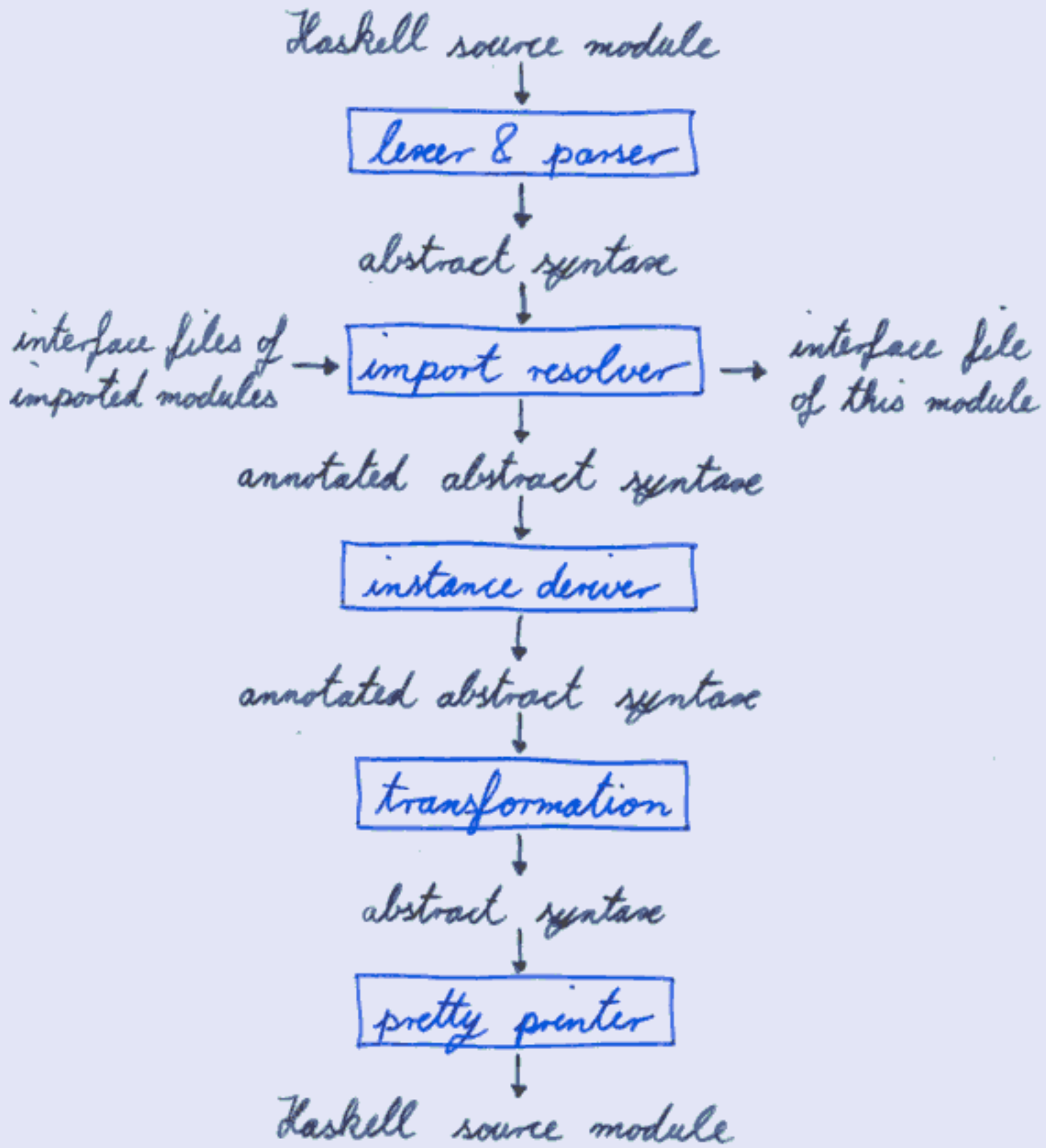
$$\text{ap2} :: \text{RefSrcPos} \rightarrow \text{RefExp} \rightarrow R (\text{Fun } a (\text{Fun } b \ z)) \rightarrow R a \rightarrow R b \rightarrow R z$$

$$\text{fun1} :: \text{RefAtom} \rightarrow \text{RefSrcPos} \rightarrow \text{RefExp} \rightarrow (R a \rightarrow \text{RefExp} \rightarrow R z) \rightarrow R (\text{Fun } a \ z)$$

$$\begin{aligned} \text{fun2} :: \text{RefAtom} \rightarrow \text{RefSrcPos} \rightarrow \text{RefExp} \rightarrow (R a \rightarrow R b \rightarrow \text{RefExp} \rightarrow R z) \\ \rightarrow R (\text{Fun } a (\text{Fun } b \ z)) \end{aligned}$$

only after evaluation of function know if application is saturated

Rat-trans



Transformation

$f :: \text{Bool} \rightarrow \text{Bool}$
 $f\ x = g\ x$



$f :: \text{RefExp} \rightarrow \text{Fun Bool Bool}$
 $f\ p = R\ (\text{Fun } (\lambda x\ r \rightarrow \text{ap } r\ (g\ r)\ x))$
 $\quad (\text{mkAt } p\ "f")$

$f\ \text{True}$



case (let $p = \text{mkRoot}$

in $\text{ap } p\ (f\ p)\ (R\ \text{True } (\text{mkAt } p\ "True"))$) of
 $R\ x_ \rightarrow x$

data $R\ a = R\ a\ \text{RefExp}$

← augmented expressions

newtype $\text{Fun } a\ b = \text{Fun } (R\ a \rightarrow \text{RefExp} \rightarrow R\ b)$

$\text{ap} :: \text{RefExp} \rightarrow R\ (\text{Fun } a\ b) \rightarrow R\ a \rightarrow R\ b$

$\text{ap } p\ (R\ (\text{Fun } f)\ r\ f)\ a\ @\ (R\ -\ r\ a) =$

let $r = \text{mkAp } p\ r\ f\ r\ a$

in $R\ (\text{entRedex } r\ \text{'seq' case } f\ a\ r\ \text{of } R\ y\ r\ y \rightarrow \text{updResult } r\ y\ \text{'seq' } y)\ r$

Transformation: Types

data Point = P Integer Integer

→ data Point = P (R Integer) (R Integer)

sort :: Ord a => [a] -> [a]

→ gsort :: Ord a => RefSrcPos -> RefExp -> R (Fun (List a) (List a))

no defaulting!

Transformation: Pattern-Matching

$\text{reverse } [] = \dots$

$\text{reverse } (x:xs) = \dots$

$\text{reverse } p\ j = \text{fun1 } \text{areverse } p\ j\ \text{reverse}$

$\text{reverse } (R\ \text{Nil})\ j = \dots$

$\text{reverse } (R\ (\text{Cons } fx\ fxs))\ j = \dots$

- k and $n+k \rightsquigarrow \text{explicit } ==, \text{fromInteger}, \dots$
- irrefutable pattern $\rightsquigarrow$ local pattern binding
- guard $\rightsquigarrow$ continuation style

How does the transformation effect performance/complexity?

- sharing of constants

monomorphic restriction enables identification of pseudo constants

- tail recursion

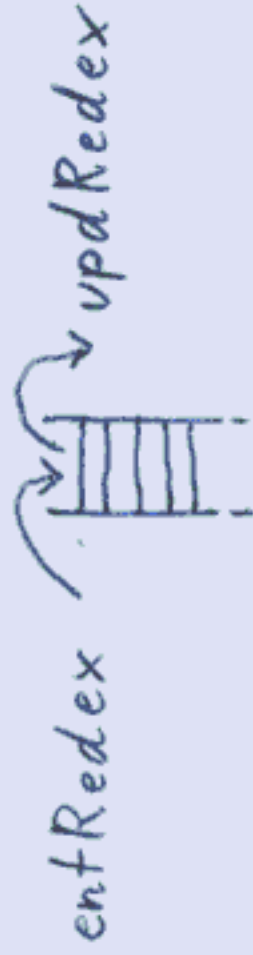
How do we handle errors?

Catching errors

- extend incomplete patterns
- define error specially
- cater Control-C and arithmetic errors with C signals
- use Haskell 10-catch
- library variants for ghc and nhc98

The Trace Stack

last entered redex without result raised error



How do we call primitive functions?

toChar :: RefExp $\rightarrow$ R Char $\rightarrow$ Char
fromChar :: RefExp $\rightarrow$ Char $\rightarrow$ R Char

toList :: (RefExp $\rightarrow$ R a $\rightarrow$ b) $\rightarrow$ RefExp $\rightarrow$ R (List a) $\rightarrow$ [b]

toString :: RefExp $\rightarrow$ R String $\rightarrow$ Prelude. String
toString = toList toChar

toFun :: (RefExp $\rightarrow$ c $\rightarrow$ R a) $\rightarrow$ (RefExp $\rightarrow$ R b $\rightarrow$ d) $\rightarrow$ RefExp
 $\rightarrow$ R (Fun a b) $\rightarrow$ (c $\rightarrow$ d)

foreign import haskell "Char.isUpper" isUpper :: Char $\rightarrow$ Bool

How do we implement trusted/untrusted code?

Wrapping does not work

- $(++) :: [a] \rightarrow [a] \rightarrow [a] \rightsquigarrow \dots \text{fromList } (\text{toList } x ++ \text{toList } y) \dots$
- $\text{elem} :: \text{Eq } a \Rightarrow a \rightarrow [a] \rightarrow \text{Bool}$
- $\text{gelem} :: \text{Eq } a \Rightarrow \text{Ref SrcPos} \rightarrow \text{Ref Exp} \rightarrow R (\text{Fun } a (\text{Fun } (\text{List } a) \text{ Bool}))$

Combinators for Trusting

- `record`: call, result, traced sub-computation
- `demand-driven recording of result`

Conclusions

- Tracing through Program Transformation
 - Hood: library + manual transformation
 - Evgia: a special compiler
 - a Haskell interpreter
- Transforming Haskell (powerful but large with irregularities)
- Compiler independent



<http://www.cs.york.ac.uk/fp/hat>